

# **DNA-based Cryptography**

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*Biotechnological Methods* (e.g., recombinant DNA) have been developed for a wide class of operations on DNA and RNA strands

*Biomolecular Computation (BMC)* makes use of such biotechnological methods for doing computation

- Uses DNA as a medium for ultra-scale computation
- Comprehensive survey of Reif [R98]
- splicing operations allow for universal computation [Head92].
- BMC solution of combinatorial search problems:

Hamiltonian path problem [Adleman94]

Data Encryption Standard (DES)  
[Boneh, et al 95] [Adleman, et al 96]

*ultimately limited* by volume requirements, which may grow exponentially with input size.

# *DNA Storage of Data*

- A medium for *ultra-compact information storage*: large amounts of data that can be stored in compact volume.
- Vastly exceeds storage capacities of conventional electronic, magnetic, optical media.
- A *gram* of DNA contains  $10^{21}$  DNA bases  
=  $10^8$  *tera-bytes*.
- A few grams of DNA may hold *all data stored in world*.
- Most recombinant DNA techniques are applied at concentrations of 5 grams of DNA per liter of water.

# *DNA Data Bases:*

- A “wet” data base of *biological data*

*natural DNA* obtained from biological sources may be recoded using nonstandard bases [Landweber,Lipton97], to allow for subsequent BMC processing.

- DNA containing data obtained from more conventional *binary storage media*.

*input and output of the DNA data* can be moved to conventional binary storage media by *DNA chip arrays*

binary data may be *encoded* in DNA strands by use of an alphabet of short oligonucleotide sequences.

## Associative Searches within DNA databases:

- methods for fast associative searches within DNA databases using hybridization [Baum95]
- [Reif95] data base join operations and various massively parallel operations on the DNA data

# *Cryptography*

Data security and cryptography are *critical* to computing data base applications.

*Plaintext*: non-encrypted form of message

*Encryption*: process of scrambling plaintext message, transforming it into an encrypted message (*cipher text*).

**Example:**

a fixed codebook provides an initial mapping from characters in the finite plaintext alphabet to a finite alphabet of codewords,

then a sophisticated algorithm depending on a key may be applied to further encrypt the message.

*Decryption*: the reverse process of transforming the encrypted message back to the original plaintext message.

*Cryptosystem*: a method for both encryption and decryption of data.

*Unbreakable cryptosystem*: one for which successful cryptanalysis is not possible.

# Our *MAIN RESULT*:

## DNA-based, molecular cryptography systems

- plaintext message data encoded in DNA strands by use of a (publicly known) alphabet of short oligonucleotide sequences.
- Based on *one-time-pads* that are in principle *unbreakable*.

## One-time-pads may be practical for DNA:

Practical applications of cryptographic systems based on one-time-pads are *limited in conventional electronic media*, by the size of the one-time-pad.

DNA provides a much more *compact storage media*, and an extremely small amount of DNA suffices even for huge one-time-pads.

## Our DNA one-time-pad encryption schemes:

- a *substitution method* using libraries of distinct pads, each of which defines a specific, randomly generated, pair-wise mapping
- an *XOR scheme* utilizing molecular computation and indexed, random key strings

# *Applications of DNA-based cryptography systems*

- the encryption of (recoded) *natural DNA*
- the encryption of DNA encoding *binary data*.

Methods for *2D data input and output*:

- use of *chip-based DNA micro-array* technology
- transform between conventional binary storage media via (photo-sensitive and/or photo-emitting) DNA chip arrays

# *DNA Steganography Systems:*

- *secretly tag* the input DNA
- then *disguise* it (without further modifications) within collections of other DNA.
- original plaintext is *not actually encrypted*
- very appealing due to *simplicity*.

## **Example:**

**DNA plaintext messages are appended with one or more secret keys**

**resulting appended DNA strands are hidden by mixing them within many other irrelevant DNA strands (e.g., randomly constructed DNA strands).**

**[Clelland, Risca, and Bancroft]**

**genomic steganography:**

**techniques using amplifiable microdots**



# Our *RESULTS* for *DNA Steganography Systems*:

- **Potential Limitations of these DNA Steganography methods:**

    Show certain DNA steganography systems can be *broken*, with some assumptions on information theoretic entropy of plaintext messages.

- We also discuss various modified DNA steganography systems which appear to have *improved security*.

# *Organization of Talk*

- ∇ *Introduction* of BMC and cryptography terminology, and results.
- ∇ *Unbreakable DNA cryptosystems* using randomly assembled *one-time pads*.
- ∇ Example of a *DNA cryptosystem for two dimensional images*, using a DNA chip for I/O and also using a randomly assembled one-time pad.
- ∇ *DNA Steganography Techniques*:
  - show that they can be *broken* with some modest assumptions on the entropy of the plaintext, even if they employ perfectly random one-time pads.
  - Provide possible improvements
- ∇ **Conclusions**

# *Cryptosystems Using Random One-Time Pads*

Use *secret codebook* to convert short segments of plaintext messages to encrypted text:

Must be *random* codebook

Codebook can be used only *once*

*In secret*, assemble a large one-time-pad in the form of a DNA strand:

randomly assembled from short oligonucleotide sequences, isolated, and cloned.

One-time-pad *shared in advance* by both the sender and receiver of the secret message:

requires initial communication of one-time-pad between sender and receiver

facilitated by compact nature of DNA

# *A DNA Cryptosystem Using Substitution*

*Substitution* one-time-pad encryption:

- a substitution method using libraries of distinct pads, each of which defines a specific, randomly generated, pair-wise mapping.
- The decryption is done by similar methods.

*Input:*

plaintext binary message of length  $n$ ,  
partitioned into plaintext words of fixed length,

*Substitution One-time-pad:*

a table randomly mapping all possible strings of plaintext words into cipher words of fixed length, such that there is a unique reverse mapping.

*Encryption:*

by substituting each  $i$ th block of the plaintext with the cipher word given by the table, and is decrypted by reversing these substitutions.

# *DNA Implementation of Substitution One-time-pad Encryption:*

- *plaintext* messages:
  - one test tube of short DNA strands
- *encrypted* messages:
  - another test tube of different short DNA strands

Encryption by *substitution*:

maps these in a random yet reversible way

plaintext is converted to cipher strands and plaintext strands are removed

DNA *Substitution one-time pads*:

use long DNA pads containing many segments:  
each segment contains a cipher word followed by a plaintext word.

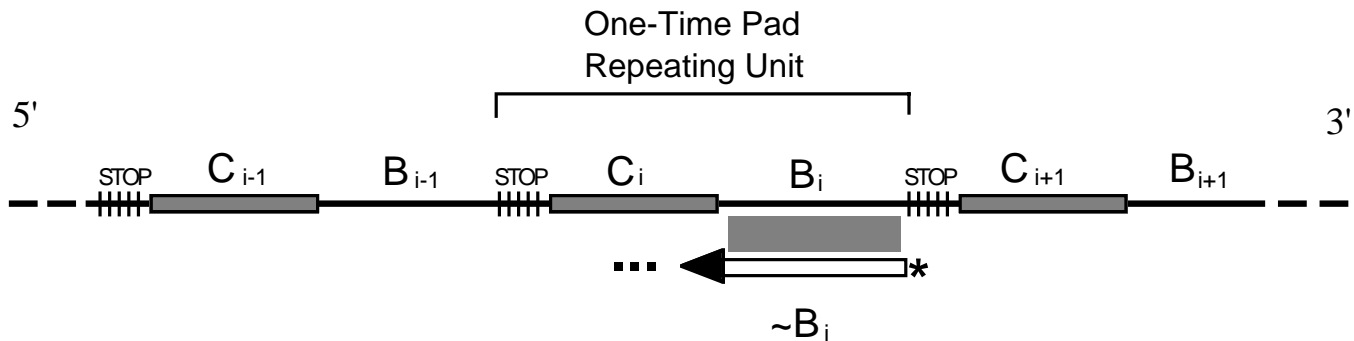
cipher word: acts as a hybridization site for binding of a primer

cipher word is appended with a plaintext word to produce word-pairs.

These word-pair DNA strands used as a lookup table in conversion of plaintext into cipher text.

# *One-time-pad DNA Sequence:*

- Length  $n$
- Contains  $d = n/(L_1 + L_2 + L_3)$  copies of repeating unit:



*Repeating unit* made up of:

- $B_i$  = a cipher word of length  $L_1 = c_1 \log n$
- $C_i$  = a plaintext word length  $L_2 = c_2 \log n$   
Each sequence pair uniquely associates a plaintext word with a cipher word.
- Polymerase "stopper" sequence of length  $L_3 = c_3$ ,

To generate a set of oligonucleotides corresponding to the plaintext/cipher word-pair strands:

- $\sim B_i$  used as polymerase primer
- *extended* with polymerase by specific attachment of plaintext word  $C_i$ .
- *Stopper sequence* prohibits extension of growing DNA strand beyond boundary of paired plaintext word.

# **Word-pair strands are essentially: *a lookup table for a random codebook.***

**Feasibility depends upon:**

- **size of the lexicon;**
- **number of possible pads available;**
- **size, complexity, and frequency of message transmissions.**

<u>Parameter</u>	<u>Range</u>
<b>Lexicon size</b>	<b>10,000 – 250,000 words</b>
<b>Word size</b>	<b>8 – 24 bases</b>
<b>Message size</b>	<b>5 – 30% of lexicon size</b>
<b>Pad diversity</b>	<b><math>10^6 - 10^8</math></b>

**Pad diversity: total number of random pads generated during a single pad construction experiment.**

# *Codebook Libraries:*

- previous gene library construction projects [LK93, LB97]
- used in DNA word encoding methods used in DNA computation [DMGFS96, DMGFS98, DMRGF+97, FTCSC97, GDNMF97, GFBCL+96, HGL98, M96].

Use *two distinct lexicons* of sequence words:

- for cipher words
- for plaintext words.

Can *generate lexicons* by normal DNA synthesis methods:

- utilize sequence *randomization at specific positions* in sequence words.

**Example:**

**For  $N = A+C+G+T$ ,  $R = A+G$ , and  $Y = C+T$ ,**

**RNNYRNRRYN**

**produces  $2 \times 4 \times 4 \times 2 \times 2 \times 4 \times 2 \times 2 \times 2 \times 4 = 16,384$  possible sequences.**



# *Methods for Construction of DNA one-time pads.*

(1) *Random assembly* of one-time pads in solution (e.g. on a synthesis column).

- *Difficult to achieve both full coverage* and yet still avoiding possible *conflicts by repetition* of plaintext and/or cipher words.
- can set  $c_1$  and  $c_2$  large so probability of repeated words on pad of length  $n$  is small, but coverage is be reduced.

(2) Use of *DNA chip technology* for random assembly of one-time pads

## *Advantages:*

currently commercially available (Affymetrix) chemical methods for construction of custom variants are well developed.

direct control of coverage and repetitions

# ***DNA chip Method* for Construction of DNA one-time pads.**

- an array of immobilized DNA strands,
- multiple copies of a single sequence are grouped together in a microscopic pixel.
- optically addressable
- known technology for synthesis of distinct DNA sequences at each (optically addressable) site of the array.
- combinatorial synthesis conducted in parallel at thousands of locations:

For preparation of oligonucleotides of length  $L$ , the  $4^L$  sequences are synthesized in  $4n$  chemical reactions.

## **Examples:**

- 65,000 sequences of length 8 use 32 synthesis cycles
- $1.67 \times 10^7$  sequences of length 10 use 48 cycles

**DNA Chip Method for**

# *Construction of DNA One-time pads*

- **plaintext and cipher pairs constructed:**
- **nearly complete coverage of the lexicon on each pad, nearly unique word mapping between plaintext and cipher pairs.**
- **resulting cipher word, plaintext word pairs can be assembled together in random order (with possible repetitions) on a long DNA strand by a number of known methods:**
  - blunt end ligation**
  - hybridization assembly with complemented pairs [Adleman97]**
- **Cloning or PCR used to amplify the resulting one-time pad.**

# *XOR One-time-pad* (Vernam Cipher) Cryptosystem

*One-time-pad S:*

a sequence of independently distributed random bits

**M:** a plaintext binary message of n bits

- *Encrypted bits:*

$$C_i = M_i \text{ XOR } S_i \text{ for } i = 1, \dots, n.$$

**XOR:** given two Boolean inputs, yields 0 if the inputs are the same, and otherwise is 1.

- *Decrypted bits:*

Use commutative property of XOR

$$\begin{aligned} C_i \text{ XOR } S_i &= (M_i \text{ XOR } S_i) \text{ XOR } S_i \\ &= M_i \text{ XOR } (S_i \text{ XOR } S_i) \\ &= M_i. \end{aligned}$$

# *DNA Implementation of XOR One-time-pad Cryptosystem*

- *plaintext messages:*  
one test tube of short DNA strands
- *encrypted messages:*  
another test tube of different short DNA strands

## *Encryption by XOR One-time-pad:*

maps these in a random yet reversible way  
plaintext is converted to cipher strands and  
plaintext strands are removed

For *efficient* DNA encoding:

use *modular base 4* (DNA has four nucleotides)

**Encryption:**

addition of one-time-pad elements modulo 4

**Decryption:**

subtract one-time-pad elements modulo 4

# *Details of DNA Implementation of XOR One-time-pad Cryptosystem*

- Each plaintext message has appended unique *prefix index tag* of length  $L_0$  indexing it.
- Each of one-time-pad DNA sequence has appended unique *prefix index tag* of same length  $L_0$ , forming *complements* of plaintext message tags.
- Use Recombinant DNA techniques (annealing and ligation) to *concatenate into a single DNA strand* each corresponding pair of a plaintext message and a one-time-pad sequence
- These are *encyphered by bit-wise XOR computation*:  
fragments of the plaintext are converted to cipher strands using the one-time-pad DNA sequences, and  
plaintext strands are removed.

*Reverse decryption* is similar:

use commutative property of bit-wise XOR operation.

# ***BMC Methods to effect bit-wise XOR on Vectors.***

Can adapt BMC methods for *binary addition*:

- similar to bit-wise XOR computation
- can *disable carry-sums* logic to do XOR

## ***BMC techniques for Integer Addition:***

- (1) [Guarnieri, Fliss, and Bancroft 96] first BMC addition operations (on single bits).
- (2) [Rubin et al 98, OGB97,LKSR97,GPZ97] permit chaining on n bits.
- (3) Addition by *Self Assembly* of DNA tiles  
[Reif,97][LaBean, et al,99]





# ***XOR by Self Assembly of DNA tiles***

[1] For each bit  $M_i$  of the message, construct sequence  $a_i$  that represents the  $i$ th bit.

[2] Scaffold strands for binary inputs to the XOR:

- Using linkers, assemble message  $M$ 's  $n$  bits into scaffold strand sequence  $a_1 a_2 \dots a_n$ ,
- One-time-pad is further portion scaffold strand  $a'_1 a'_2 \dots a'_n$  is created from random inputs

[3] add output tiles; annealing give self assembly of the tiling.

[4] adding ligase yeilds reporter strand

$R = a_1 a_2 \dots a_n \cdot a'_1 a'_2 \dots a'_n \cdot b_1 b_2 \dots b_n$   
where  $b_i = a_i \text{ XOR } a'_i$ , for  $i = 1, \dots, n$ .

[5] reporter strand is extracted by melting away the tiles' smaller sequences, and purifying.

contains concatenation of:

input message, encryption key, ciphertext

[6] Using a marker sequence:

ciphertext can be excised and separated based on its length being half that of remaining sequence.

[7] Ciphertext can be stored in a compact form

# *DNA Cryptosystem for 2D Images*

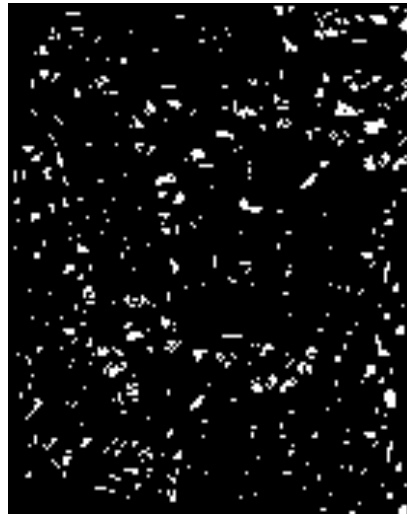
using:

- DNA Chip
- Randomly Assembled One-Time Pad

Encryption and Decryption of 2D images recorded on microscopic arrays of a DNA chip:



*Message*



*Encrypted*



*Decrypted*

Simulated patterns observed by fluorescence microscopy of the DNA I/O chip.

DNA Cryptosystem consists of:

- Data set to be encrypted: *2-dimensional image*
- *DNA Chip* bearing immobilized DNA strands:
  - contains an addressable array of nucleotide sequences immobilized s.t. multiple copies of single sequence grouped together in a microscopic pixel.
- *Library of one-time pads* encoded on long DNA strand

# *Initialization and Message Input*

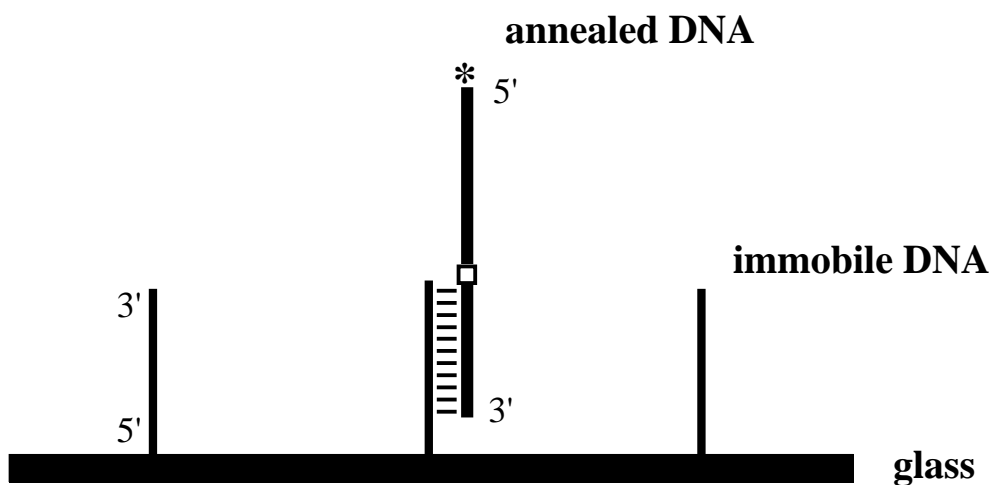
- **Fluorescent-labeled, word-pair DNA strands are prepared from a substitution pad codebook**
- **These are annealed specifically to their sequence complements at unique sites (pixels) on the DNA chip.**
- **The message information is transferred to a photo mask with transparent (white) and opaque (black) regions:**



*Message Input to DNA Chip*

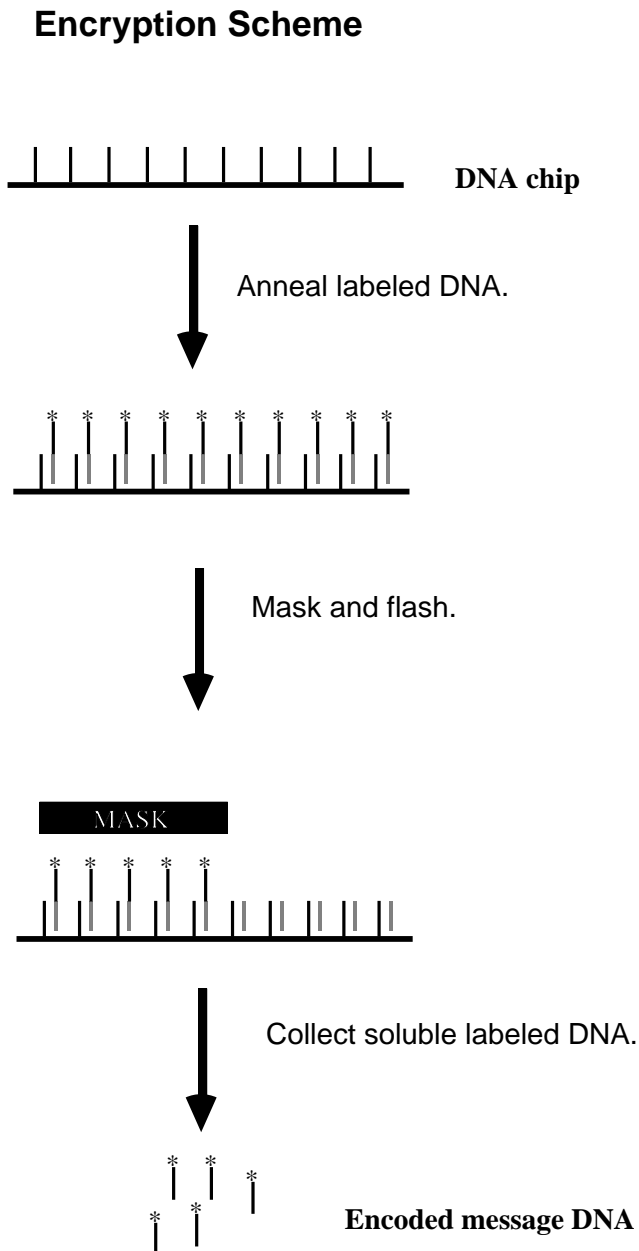
# Initialization and Message Input

- **Immobile DNA strands** are located on the glass substrate of the chip in a sequence addressable grid.
- **Word-pair strands** are prepared from a random substitution pad:
  - the 5' (*unannealed*) end carries a cipher word
  - the 3' (*annealed*) end carries a plaintext word.
  - contain a *photo-cleavable* base analog between two sequence words (added to 3' end of cipher word during oligo synthesis)



- The *annealed DNA* contains:
  - a fluorescent label on its 5' end (asterisk);
  - a codebook-matching sequence word (not base-paired on the chip);
  - a photo-labile base (white square) capable of cleaving the DNA backbone; and
  - a chip-matching word (base-paired to immobile strand).

# Encryption Procedure:



[1] start with DNA chip displaying sequences *complementary* to plaintext lexicon.

[2] fluorescent-labeled word-pair strands from one-time-pad are *annealed* to chip at pixel bearing complement to plaintext 3' end.

[3] mask protects some pixels from a light-flash. At unprotected regions, DNA is *cleaved* between plaintext and cipher words.

[4] cipher word strands, still labeled with fluorophore at 5' ends, are collected and transmitted as *encrypted message*.

# *Encryption of the Message*

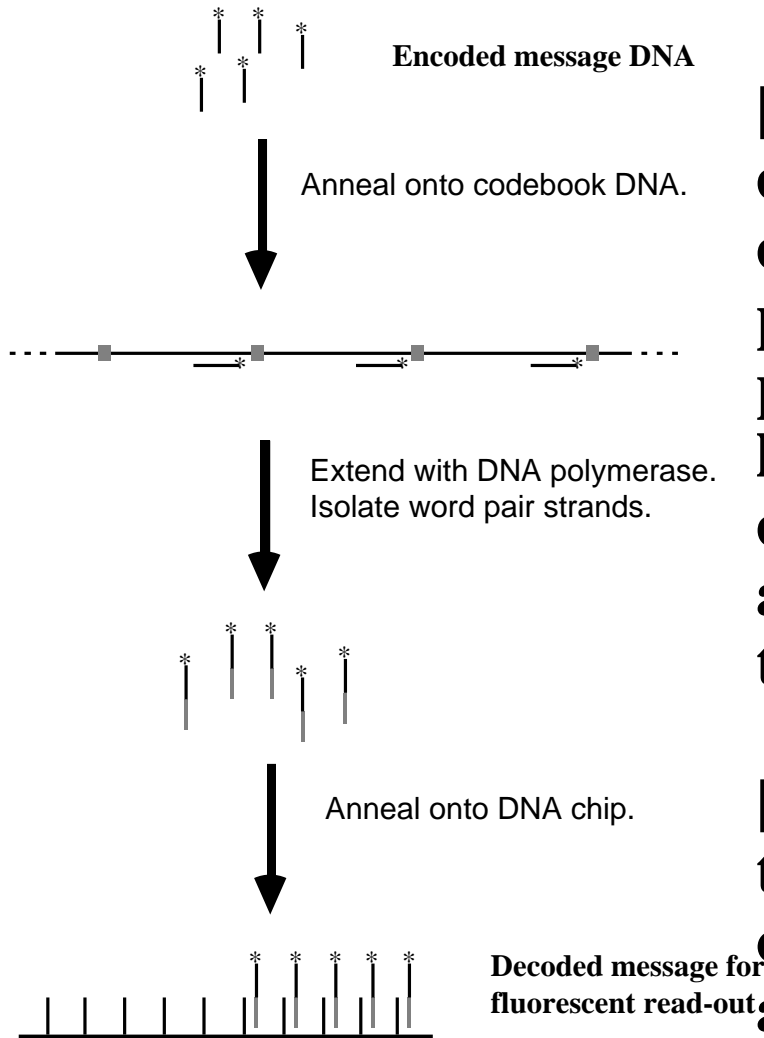
- Following a light-flash of mask-protected chip, *annealed* oligonucleotides beneath transparent mask pixels are *cleaved* at a photo-labile position: their 5' sections are dissociated from annealed 3' section and collected in solution.
- This test tube of strands is *encrypted message*.
- Annealed oligos *beneath opaque mask* are unaffected by light-flash and can be *washed off chip*.
- If encrypted message oligos are reannealed onto a (washed) DNA chip, message information would be *unreadable*:



*Simulated Read-Out of Encrypted Message from DNA Chip*

# Decryption Procedure:

## Decryption Scheme



[1] word-pair strands constructed, *appending* cipher word with proper plaintext word, by polymerase extension or lop-sided PCR using cipher words as primer and one-time-pad as template.

[2] cipher strands *bind* to their specific locations on the pad and are appended with their plaintext partner.

[3] binding reformed *word-pair strands* to DNA chip and reading message by fluorescent microscopy.

# *Decryption of the Message*

- use the fluorescent labeled oligos as primers in one-way (lopsided) PCR with the same one-time codebook which was used to prepare the initial word-pair oligos.
- When word-pair PCR product is bound to the same DNA chip, the decrypted message is revealed:



**Decrypted Message**

*Simulated Read-Out of Decrypted Message from DNA Chips*



# *Steganography*

**a class of techniques that hide secret messages within other messages:**

**plaintext is not actually encrypted but is instead disguised or hidden within other data.**

## *Historical examples:*

- **use of grills that mask out all of an image except the secret message,**
- **micro-photographs placed within larger images**
- **invisible inks, etc.**

## *Disadvantages:*

- **Cryptography literature generally consider conventional steganography methods to have *low security*:**

**steganography methods have been often broken in practice [Kahn67] and [Schneier96]**

## *Advantages:*

- **it is very appealing due to its *simplicity*.**

# *DNA Steganography Techniques:*

- take one or more input DNA strands (considered to be the plaintext message)
- append to them one or more randomly constructed “*secret key*” strands.
- Resulting “*tagged plaintext*” DNA strands are *hidden* by mixing them within many other additional “distracter” DNA strands which might also be constructed by random assembly.

## *Decryption:*

- Given knowledge of the “secret key” strands,
- Resolution of DNA strands can be decrypted by a number of possible known recombinant DNA separation methods:

plaintext message strands may be separated out by hybridization with the complements of the “secret key” strands might be placed in solid support on magnetic beads or on a prepared surface.

These separation steps may combined with amplification steps and/or PCR

# *Cryptanalysis of DNA Steganography Systems:*

DNA steganography system's security is entirely dependent on degree that message DNA strands are *indistinguishable* from “distracter” DNA strands.

## *Cryptanalysis Assumptions:*

- no knowledge of the “secret key” strands
- secret tags are indistinguishable from “distracter” DNA strands.
- plaintext is not initially compressed, and comes from a source (e.g., English or natural DNA) with Shannon information theoretic entropy  $E_s > 1$
- the “distracter” DNA strands are constructed by random assembly

Then:

the original plaintext portion of “tagged plaintext” DNA strands are *distinguishable* from “distracter” DNA strands, and

the DNA Steganography System can be *broken*

# *Shannon (information theoretic)*

## *Entropy $E_s$*

- provides a measure of the factor that a source can be *compressed* without loss of information.

### Examples:

many images have entropy nearly 4

English text has entropy about 3

computer programs have entropy about 5

most DNA have entropy range 1.2 to 2

### *Lossless Data Compression [Lempel-Ziv 77]*

*Input:* text string of length  $n$  with entropy  $E_s$

[1] Form a *dictionary*  $D$  of the  $d = n/L$  most frequently occurring subsequences of length at least  $L = E_s \log_2 n$  in the known source distribution.

[2] In place of subsequences of the input text matching with elements of the dictionary  $D$ , *substitute their indices* in the dictionary  $D$ .

# *Cryptanalysis of DNA Steganography Systems:*

***Input:*** test tube T containing:

a mixture of “tagged plaintext” DNA strands mixed with a high concentration of “distracter” DNA strands, of length  $n$ .

- form a *dictionary* D of the  $d = n/L$  most frequently occurring subsequences of length at least  $L = E_s \log_2 n$  in the known plaintext source distribution.
- Give procedure for *separating* out plaintext message strands by repeated rounds of hybridization with complements of elements of D.

$r(T) =$  *ratio of concentration* of “distracter” DNA strands to “tagged plaintext” DNA strands.

On each *round of separation*:

form a new test tube  $F(T)$  with expected  $r(F(T))$  considerably *reduced* from the previous ratio  $r(T)$ .

## *Separation Procedure:*

[1] Pour a fraction  $s = 1/2$  of volume of current test tube  $T$  into a test tube  $T_1$  and pour remaining fraction  $1-s$  of  $T$  into test tube  $T_2$ .

[2] Choose a *random text phrase*  $x$  in  $D$  (not previously considered in a prior trial), and using Watson-Crick complement of  $x$ , do a *separation* on test tube  $T_2$ , yielding a new test tube  $T_3$  whose contents are only DNA strands containing phrase  $x$ .

[3] Pour contents of test tubes  $T_1$  and  $T_3$  into a new test tube  $F(T)$ .

- Ratio  $r(F(T))$  of “distracter” DNA strands to plaintext DNA will expect to *decrease* from original ratio  $r(T)$  by a constant factor  $c < 1$
- After  $O(\log(r/r'))$  repeated rounds of this process, ratio of concentration in test tube  $T$  will expect to *decrease* from initially  $r = r(T)$  to any given smaller ratio  $r'$ .

# *Another cryptanalysis technique for breaking steganographic systems:*

Cryptanalysis using “*hints*” that disambiguate plaintext.

**Example:**

- wish to make secret the DNA of an individual (e.g., the President)
- use an improved steganography system where “distracter” DNA strands (that are mixed with DNA of an individual) are DNA from a similar but not identical genetic pool.

steganography system may often be *broken* by use of distinguishing “*hints*” concerning DNA of the individual

e.g., the individual might have a particular set of observable expressed gene sequences (e.g., for baldness, etc.).

**These hints may allow for subsequent identification of the full secret DNA:**

use of a series of separation steps with complement of portions of known gene sequences.

# *Improved DNA Steganography Systems with Enhanced security:*

**Idea:** make it more difficult to *distinguish* probability distribution of plaintext source from that of “distracter” DNA strands.

## *(1) Mimicking Distribution of “Distracter” DNA:*

- use improved construction of the set of “distracter” DNA strands, so distribution better mimics the plaintext source distribution
- construct the “distracter” DNA strands by random assembly from elements of Lempel-Ziv dictionary.
- *Drawback:* Cryptanalysis using “hints” that disambiguate plaintext.

## *(2) Compression of Plaintext.*

- recode the plaintext using a universal lossless compression algorithm (e.g., Lempel-Ziv 77).
- resulting distribution of the recoded plaintext approximates a universal distribution, so uniformly random assembled distracter sequences may suffice to provide improved security.
- *Drawback:* unlike conventional steganography methods, plaintext messages need to be preprocessed.



# *Conclusion and Open Problems*

Presented an initial investigation of DNA-based methods for Cryptosystems.

- *Main Results* for DNA one-time-pads cryptosystems:

Gave DNA substitution and XOR methods based on one-time-pads that are in principle *unbreakable*.

Gave an implementation of our DNA cyptography methods including *2D input/output*.

- **Further Results for *DNA Steganography*:**

a certain class of DNA steganography methods offer only limited security; can be *broken* with some reasonable assumptions on entropy of plaintext messages.

modified DNA steganography systems may have *improved security*.

## *Open Problem:*

Show whether DNA steganography systems with natural DNA plaintext input can or cannot be made to be *unbreakable*.