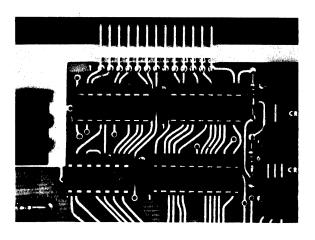
# **USER'S MANUAL**

# The Speakeasy<sup>tm</sup> cassette/general purpose interface



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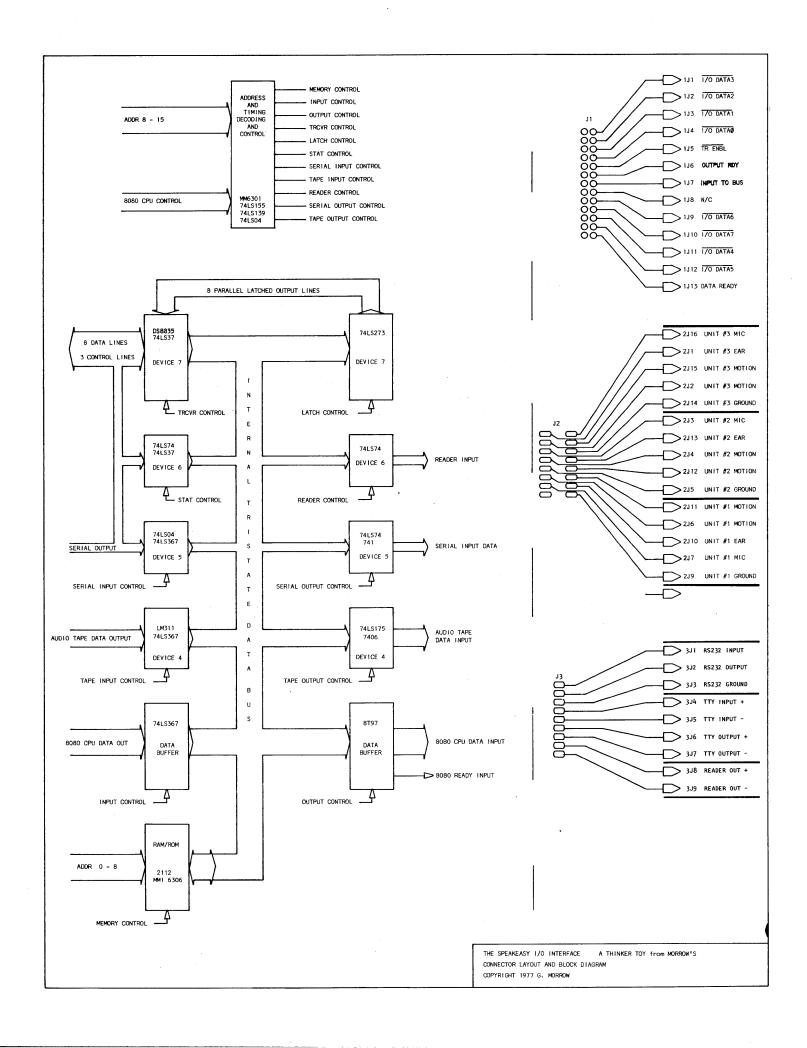


# THE SPEAKEASY<sup>TM</sup> CASSETTE/GENERAL PURPOSE INTERFACE

#### INTRODUCTION

The Speakeasy was the first and is still the most versatile of the "intelligent" S-100 interface boards. It pioneered the use of RAM and ROM and PROM on the I/O board itself. There are three completely separate interfaces on the Speakeasy board. Yet, the board is not dense. There is very little chance of creating solder bridges because there is sufficient clearance for all the parts. Other I/O boards which are as versatile as the Speakeasy have upwards to seventy integrated circuits and traces ten mills wide. The reason for the vast difference in circuit complexity is that the Speakeasy has its complexity concentrated in its PROMs. The intelligence of the CPU is harnessed to keep the circuit layouts simple and clean.

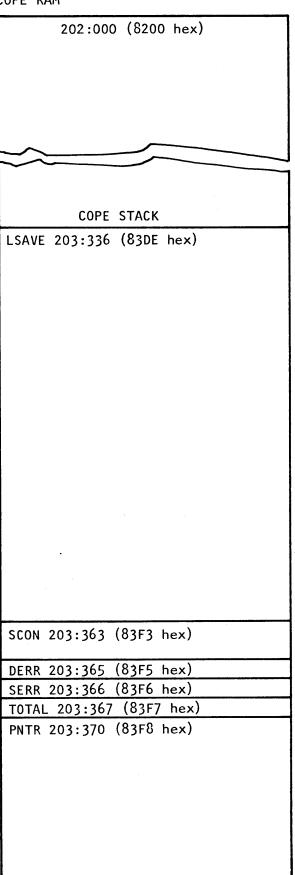
However, there are certain trade-offs incurred by this novel approach to the I/O interfacing: the CPU communicates with the Speakeasy interfaces much differently than the standard "dumb" I/O interfaces. The purpose of this manual is to ease the problems associated with this novel and more efficient method of controlling I/O devices. The staff at Morrow's has worked diligently to make this manual as clear and complete as possible and, most importantly, useable. Comments on improving the material will be welcomed. Errors are always difficult to completely eliminate from a technical document; your help in finding them is solicited. Please feel free to write us with your comments.



#### COPE ROM

FLUX	200:000	(8000	hex)	
WTAPE	200:034	(801C	hex)	
TREAD	200:115	(804D	hex)	
COPE 2	01:012 (	810A I	nex)	
CHECK	201:115	(814D	hex)	
BOOTS	201:137	(815F	hex)	
INPUT	201:157	(816F	hex)	
DELAY	201:240	(81AO	hex)	
SROUT	201:263	(81B3	hex)	
DETCT	201:313	(81CB	hex)	

#### COPE RAM



#### PRINCIPLES OF OPERATION

## THE TAPE CASSETTE INTERFACE

#### GENERAL

The I/O board contains interface circuitry that allows an S-100 compatible computer to communicate with three audio cassette player/recorders. The interface can read from any one of the three tape channels and can write on any combination of the three tape channels. The data standard that the interface uses is the "Kansas City" 300 baud data format agreed upon several years ago by a diverse group of manufacturers of personal computing equipment.

### THE "KANSAS CITY" DATA STANDARD

The Kansas City Standard (KCS) is a data format which specifies how digital data shall be encoded into tones that can be recorded on low cost audio cassette/player recorders. KCS defines a logic 1 to be a 2400 hertz signal which is eight complete cycles long. A logic 0 is a 1200 hertz signal which is four complete cycles long. KCS also specifies how logic 0s and 1s shall be grouped together in bytes on the audio cassette: there is to be at least two logic 1 bits (rest bits) followed by a logic 0 bit (start bit) followed by eight data bits.

The KCS is a conservative data format. There is a large factor of redundancy (a factor of eight for logic 1 and a factor of four for logic 0) which means that this is an extremely reliable method of data storage and retrieval. As with most things, the reliability of the KCS is not without penalties. In this case, the penalty is the relatively slow rate that data is stored or retrieved -- 300 bits per second. However, the first concern of any decent standard should be data integrity and in this sense the KCS performs admirably.

#### SUITABLE AUDIO CASSETTE PLAYER/RECORDERS

There are a wide variety of audio cassette units available which are suitable for use with the I/O board. However, there are several features which a cassette unit should have.

- (1) The player should have a digital counter so that file positions on tape can be easily found.
- (2) The unit should have AGC (automatic gain control) for making recordings so that volume settings can be restricted to one value rather than requiring one adjustment for writing and another for reading.
- (3) There should be an input jack that allows a remote switch to activate the unit.
- (4) There should be an auxiliary input other than the microphone input which will accommodate a radio or record player pre-amp output.

The signal level from the cassette interface of the I/O board is approximately 4 volts peak-to-peak and is too high for microphone input. This high level signal has a much better signal-to-noise ratio than the 50 millevolt signal necessary for the microphone input. During recording, the microphone input must be disabled. Sometimes, the microphone is disabled when a jack is inserted in the auxilliary input but more often the cassette unit comes supplied with a small plastic plug that is inserted into the microphone jack to disable the input.

#### THE SOFTWARE INTERFACE

A unique design feature of the I/O board is the use of PROM and RAM to link the 8080 CPU intimately with the cassette player/recorders connected to the board. Through the use of software in PROM, the I/O board directs the CPU to measure precise time intervals with time delay subroutines and to use these timing routines to generate the 1200/2400 Hz wave forms that are sent out to the recorder. Using other software routines stored in the on-board PROM, the CPU measures the time between zero-crossings of the incoming signals from the cassette player to convert the 1200/2400 Hz wave forms into digital information.

The software in the PROMs in highly modularized. For input, specific routines measure the time between wave form zero-crossings. Other routines use the zero-crossing measurements to build individual bits while other software uses the bit-building software to build bytes. High level routines take assembled bytes and transfer them to memory.

For output, the high level routines transfer bytes to routines which disassemble the bytes to bits. The bits are in turn transferred to subroutines which convert them to 1200/2400 Hz wave forms through the use of time measurement software.

The software has been designed to read or write entire blocks of information to or from memory with no intervention from external software. A program such as an editor or BASIC communicates with software of the tape cassette interface by issuing a CALL COPE instruction. Certain CPU registers must be initialized to furnish detailed information to the I/O board software about the number of bytes to be transferred, the address of the first byte and the type of transfer (read, write, verify). The rest of this section explains the details of how to use the powerful COPE software which is in the PROMs of the I/O board.

#### COPE

COPE is an acronym for Cassette OPerating Executive. It is the name given to that body of software in the PROMs which effects data transfers to and from the cassette player/recorders connected to the I/O board. Below are the kinds of tasks COPE will perform:

- (1) Read a block of data (from 1 to 65k bytes) from any one of three cassette players.
- (2) Verify a block of data on a cassette player against a block of data in memory.

- (3) Move across a block of data on a cassette player.
- (4) Write a block of data from memory to any one or combination of three cassette recorders.

#### COPE COMMAND STRUCTURE

The COPE software has exactly <u>one</u> entry point: 201:012 octal or 810A hex and is almost without exception entered via a CALL instruction (315 octal, CD hex). How does the COPE software distinguish between reading, writing, etc., if there is only one entry point? The A register or "accumulator" of the 8080A serves as the primary command register and the C register acts as a secondary command register. That is, the patterns or values present in the A and C registers when a "CALL COPE" instruction is executed determines what kind of operation the COPE software will perform. The pattern in the C register is ignored if the value in the A register specifies a write command. If the A register does not specify a write command, the value of the C register is used to specify whether the command is a read, verify or move.

#### A REGISTER COMMAND FORMATS

The convention used below to specify bit positions in 8080 registers is the same as the one used in the 8080 User's Manual published by INTEL CORP: the higher the bit number, the greater the significance the bit has in forming a value of the register.

A register command structure:

bit 7 bit	6 bit 5	bit 4 bit	3 bit 2	bit 1	bit Ø

- Bit Ø The write/not write bit. If this bit is a one, data will be read from the computer's memory and written on any combination of the three cassette recorders connected to the interface. If this bit is a zero, data from exactly one of the cassette players is to be: (i) read from the tape into the computer's memory; (ii) read from the tape and compared with memory; or (iii) simply passed over so as to access the next block of data on the tape. The C register is used as a secondary command register when bit O is a zero to decide which of the three alternatives should be performed by the COPE software.
- <u>Bit 1</u> Cassette channel #1 bit. This bit controls the motion of the cassette player/recorder connected to channel #1 of the I/O board. A one in this bit position turns the motor and electronics "on" while a zero turns them "off."
- <u>Bit 2</u> Cassette channel #2 bit. This bit controls the motion of the cassette player/recorder connected to channel #2 of the I/O board. A one in this bit position turns the motor and electronics "on" while a zero turns them "off."
- <u>Bit 3</u> Cassette channel #3 bit. This bit controls the motion of the cassette player/recorder connected to channel #3 of the I/O board. A one in this bit position turns the motor and electronics "on" while a zero turns them "off."

Principles of Operation

Bit 4 - Always a zero.

Bit 5 - Always a zero.

Bit 6 - Always a zero.

 $\underline{\text{Bit 7}}$  - If this bit is a one,  $\underline{\text{all}}$  the motion control relays will be turned off when COPE completes a command. There are exceptions, however, and one will be described later.

#### C REGISTER COMMAND FORMATS

The value or pattern of the C register is of no importance to the COPE software unless bit 0 of the A register is zero. When bit 0 of the A register is a zero, the C register has the following command structure:

bit 7 bit 6 bit 5 bit 4 bit 3 bit 2	bit 1 bit 0
-------------------------------------	-------------

<u>Bit 0</u> - The MOVE bit. If this bit is a one and bit 0 of the A register is a zero, no data will be transferred from tape to memory. Instead, the active tape transport will move the tape across a block of data. The length of this block is discussed in the next section.

Bit 1 - Always a zero.

Bit 2 - Always a zero.

Bit 3 - Always a zero.

Bit 4 - Always a zero.

Bit 5 - Always a zero.

Bit 6 - The VERIFY bit. If this bit is a one and bit 0 of the A register is a zero, a "compare" or "verify" operation will be performed. Data will be read from the transport connected to the active channel and compared bytefor-byte with memory data. If any data on the tape should differ from its counterpart in memory, an error flag will be set. This is discussed later.

#### Bit 7 - Always a zero.

If all the bits of the C register are zero and bit 0 of the A register is zero, a read operation will be performed. The COPE software will transfer a block of data from the cassette player connected to the active channel and write this data into the computer's memory.

#### COPE DATA BLOCK SPECIFICATIONS

As mentioned above, COPE software performs operations on blocks of data. When a block resides in a computer's memory, it has three attributes:

- (1) A starting address.
- (2) An ending address.
- (3) A length (in bytes).

These attributes are not independent. Given any two, the third can be calculated. The two that COPE uses are the starting address of the data block and its length. COPE requires the starting address to be in the H-L register pair and the block length to be in the D-E register pair.

When COPE writes data from memory onto one or more cassette channels, five CPU registers must be initialized before transferring control to the COPE software:

- (1) The A register must be initialized with the proper pattern.
- (2) The H-L register pair must be initialized with the starting address of the data block that will be written on the tape(s).
- (3) The D-E register pair must be initialized with the length of the data block.

When COPE reads or verifies data on tape, six registers must be initialized before transferring control to the COPE software:

- (1) The A register must be initialized with the proper pattern.
- (2) The C register must specify a read or verify operation.
- (3) The H-L register pair must be initialized with the starting address of the data block.
- (4) The D-E register pair must be initialized with the length of the data block.

When COPE spaces the tape head across a block of data, four registers must be initialized before transferring control to the COPE software.

- (1) The A register must be initialized with the proper values.
- (2) The C register must specify a move operation.
- (3) The D-E register pair must be initialized with the length of the data block.

The H-L register pair does not have to be initialized since no data is transferred during a move operation. However, the value in this register pair is not preserved during a COPE move operation. In fact, any COPE operation changes the values of <u>all</u> the CPU registers <u>except</u> the stack pointer.

There are two memory locations in the RAM of the I/O board that COPE uses for error reporting. DERR (203:365 octal, 83F5 hex) is the Data ERRor memory cell. This location is over-written with some non-zero value whenever an error is encountered during verify operations. COPE never initializes this cell and writes in this cell only during verify operations when an error is encountered. If, for example, five errors were encountered during a

particular verify operation, COPE would write into the DERR location five times -- each time with some non-zero value. The proper way to use the DERR location is to initialize it to zero some time prior to a verify operation. After the verify is complete, this cell should be tested. If it is still zero, no error was encountered during the verify. If it has changed from zero, an error occurred.

The other error reporting memory location is SERR (203:366 octal, 83F6 hex). This is the Status ERRor memory cell. Whenever COPE is called on to perform a READ, VERIFY, or MOVE operation, one of the first things COPE does is turn on the motor and electronics of the tape cassette player that is connected to the active channel. If through an error or oversight, the cassette player is not plugged in, turned on, or connected to the channel, the COPE wave form detection software will see only spurious signals and not regular waveforms. If this situation persists for more than 45 seconds at a time, COPE will abort the operation and write some non-zero value in SERR and then return control to the calling program. As with the DERR cell, SERR should be initialized to zero before initiating a READ, VERIFY or MOVE operation when it is necessary to detect the absence of valid waveforms over an extended period of time. If a read operation were inadvertantly started before the play button of the player were pressed, no harm would be done as long as it is depressed within 45 seconds after COPE has taken control.

#### **EXAMPLES USING COPE**

This section presents several practical examples that illustrate how the technical specifications of the previous sections are used.

Example 1: Write  $(728)_{10}$  bytes of data starting at location  $(120)_{10}$  of page 7 in memory onto tape channels 2 and 3.

8080 machine code is usually expressed in octal or hex so 728 and 120 need to be converted.

$$(728)_{10} = (208)_{16} = (2:330)_{8}$$

and 
$$(120)_{10} = (78)_{16} = (178)_{8}$$

2:330 is the so-called "split octal" representation. Numbers larger than 255 must use at least two bytes of storage. 330 is the value that occupies the lower of the two bytes while 2 is the value in the high order byte. The method of conversion is as follows: calculate how many times 256 will divide the number. This is the octal number to the left of the colon. The remainder is the octal number to the right of the colon. In the present case, 256 divides 728 twice with a remainder of 216. 216 has an octal representation of 330 (3.64 + 3.8 + 0.1).

The D-E register pair must be loaded with the length of data block which is 2:330 in octal or 208 in hex. The H-L register pair must be initialized with the starting address of the data block which has the split octal representation of 7:170 and the hex representation of 778.

Write operations do not use the secondary command register so the only other register that must be initialized is the A register -- the COPE command

register. The motion should be terminated at the end of the operation so bit 7 should be a one. For all commands, bits 6, 5, and 4 must be zero. Tape channels two and three are to be active so bits 3 and 2 should be a one. Bit 1 should be a zero because tape channel 1 is not active. Finally, bit 0 must be a one so that COPE will perform a write operation. The bit pattern of A is summarized below:

bit 7	bit 6	bit 5	bit 4	bit 3	bit 2	bit 1	bit 0
1	0	0	0 0		1	0	1

This bit pattern corresponds to the octal expression 215 and the hex expression 8D. The initialization procedure to write the above data onto the tape can be expressed in a sequence of 8080 instructions.

Hex	(		0cta	1		
3E 8D	MVI	A,8D	076 203	MVI	A,203	INITIALIZE CMD REG
11 08 02	LXI	D,208	021 330 002	LXI	D,2:330	INITIALIZE BLK LENGTH REG
21 78 07	LXI	н,778	041 178 007	LXI	н,7:178	INITIALIZE START ADDR REG
CD OA 81	CALL	COPE	315 012 201	CALL	COPE	EXECUTE COPE OPERATION

The particular order of the first three instructions is not important. The CALL COPE instruction sets the operation in motion and must occur last. A word of warning concerning CALL instructions: it is vital that the 8080 CPU stack register hold an address which points to valid RAM memory whenever a CALL instruction is executed. Subroutines usually end with a return instruction. This is a special type of branch -- the branch address is furnished by the memory locations pointed to by the stack register. Under normal circumstances a CALL instruction has stored a return address in these two locations. If the stack pointer points to ROM or non-existent memory, a CALL instruction cannot store the correct return address.

Example 2: Test the write circuitry of the I/O board. In this example, test data will be written from the RAM of the I/O board onto tape channel 1. The test data along with the addresses is presented below:

0ctal																
202:1	00	000	0	01	002	0	03	004	C	05	006	0	07			
202:1	10	010	0	11	012	0	13	014	C	115	016	0	17			
202:1	20	020	0	21	022	0	23	024	C	25	026	0	27			
202:1	30	030	0	31	032		33	034		35	036		37			
202:1	40	040	0	41	042	0	43	044	C	145	046	0	47			
202:1	50	050	0	51	052	0	53	054	C	55	056	0	57			
Hex																
8240	00	01	02	03	04	05		•	08	09	0 <sup>'</sup> A	0B	00	OD	0E	0F
8250	10	11	12	13	14	15	16	17	18	19	1A	1 B	1C	1 D	1 E	1F
8260	20	21	22	23	24	25	26	27	28	29	2A	2B	2C	2 D	2E	2F

As in example 1, bit 7 of the A register should be 1 and bits 6, 5 and 4 must be 0. Because only channel 1 will be active, bits 3 and 2 are zero while bit 1 is one. Finally, since the operation is to write data onto tape, bit 0 is also 1. Thus the A register which is the COPE command register must have the following pattern:

Ь	t 7	bit 6	bit 5	bit 4	bit 3	bit 2	bit 1	bit 0
	1	0 0		0	0	0	1	1

This corresponds to the octal number 203 or the hex number 83. Again, as in example 1, because the operation is writing data onto tape, it is not necessary to initialize register C, the secondary command register. The starting address of the test data as outlined above is 202:100 octal or 8240 hex. The data block has a length of 60 octal or 30 hex.

The set of instructions to write the test data out on channel one's tape recorder will conveniently fit at the beginning of the I/O boards RAM: 202:000 octal (8200 hex).

Octal			
202:000 202:003 202:005 202:010 202:013 202:016	061 370 202 076 203 021 060 000 041 100 202 315 012 201 303 016 202	LXI SP,202:370 MVI A,203 LXI D,60 LXI H,202:100 CALL COPE JMP 202:016	INITIALIZE STACK POINTER INITIALIZE COMMAND REG LOAD THE DATA BLOCK LENGTH LOAD STARTING ADDR OF BLOCK EXECUTE COMMAND COME TO DYNAMIC HALT
Hex 8200 8203 8205 8208 820B 820E	31 F8 82 3E 83 11 30 00 21 40 82 CD OA 81 C3 OE 82	LXI SP,82F8 MVI A,83 LXI D,30 LXI H,8240 CALL COPE JMP 820E	

Load this program into the RAM of the I/O board starting at the indicated address. Connect cables from the I/O board to a suitable audio cassette player/recorder. Before connecting the motion control cable, load a cassette tape into the unit, position the tape past the blank leader, reset the digital counter on the unit to zero and connect the motion control cable. Next, depress the play/record buttons. Start the program. The tape will begin to move and will stop after a short while. When the tape stops, the above program will have completed the CALL COPE instruction and will be executing the JMP 820E dynamic halt instruction. Halt the the computer but do not turn it off. Remove the motion control and external speaker plugs from the cassette player/recorder.

Rewind the tape to the point where the recorded data started. Press the play button. There should be a steady 2400 Hz tone that lasts approximately ten seconds. The data should follow the tone and will not last more

than several seconds. After verifying the presence of the recorded data, rewind the tape so that the head is positioned over the beginning of the leader tone. Reconnect the motion control and the external speaker cables.

Example 3: Test the read circuitry of the I/O board by reading the test data written in Example 2.

The bit pattern of the A register is exactly the same as in Example 2 except that bit 0 is a zero because the desired operation is a read and not a write. Thus, the A register has the following pattern:

bit 7	bit 6	bit 5	bit 4	bit 3	bit 2	bit 1	bit 0
1	0	0 0		0	0	. 1	0

This pattern corresponds to the octal number 202 or the hex number 82. Because the command is <u>not</u> a write operation, it is necessary to initialize the secondary command register which is the C register. This register must be initialized to zero since the desired operation is not a verify or move.

For comparison purposes, the data on the tape should be read into an area of RAM memory different than where the data of the previous example was stored. A convenient place in the I/O board RAM is between locations 202:200 and 202:257 octal. The set of instructions to read the previously written test data back into the I/O board RAM starting at location 202:200 octal is presented below:

Octal			
202:030 202:033 202:035 202:037 202:042 202:045 202:050	061 370 202 076 202 016 000 021 060 000 041 200 202 315 012 201 303 050 202	LXI SP,202:370 MVI A,202 MVI C,000 LXI D,60 LXI H,202:200 CALL COPE JMP 202:050	INITIALIZE STACK POINTER INITIALIZE COMMAND REGISTER SET UP SECONDARY COMMAND REG DATA BLOCK LENGTH DATA BLOCK STARTING ADDRESS EXECUTE READ COMMAND DYNAMIC HALT
Hex			
8218 821B 821D 821F 8222 8225 8228	31 F8 82 3E 82 0E 00 11 30 00 21 80 82 CD 0A 81 C3 28 82	LXI SP,82F8 MVI A,82 MVI C,0 LXI D,30 LXI H,8280 CALL COPE JMP 8228	

Enter this program into the RAM of the I/O board starting at the indicated address. Adjust the volume setting of the cassette player at approximately 4/5 of the full volume. Occasionally, it will be necessary to experiment with the volume setting. If there is a tone control on the cassette

Octal

unit, it should be set to maximum treble. Depress the PLAY button of the cassette and start the program. The tape will start to move and in a short while, it will stop and the COPE command will be done. Stop the computer and press the stop button of the cassette unit. Examine locations 202:200 through 202:257 octal. The data should have the following pattern:

octai																
202:200		000	00	1	002	00	3	004	00	5	006	00	7			
202:210		010	01	1	012	01	3	014	01	5	016	01	7			
202:220		020	02	1	022	02	3	024	02	5	026	02	7			
202:230		030	03	1	032	03	3	034	03	5	036	03	7			
202:240		040	04	1	042	04	3	044	04	5	046	.04	7			
202:250		050	05	1	052	05	3	054	05	5	056	05	7			
Hex																
8280	00	01	02	03	04	05	06	07	08	09	OA	ОВ	0C	OD	0E	0F
8290	10	11	12	13	14	15	16	17	18	19	1A	1B	1 C	1 D	1E	1F
82A0	20	21	22	23	24	25	26	27	28	29	2A	2B	2C	2 D	2E	2F

If the data is not correct, experiment with the volume setting on the cassette player and restart the program. If errors persist, return the Speakeasy for service.

Example 4: Verify the data on the tape against data in memory.

Be sure that the data in the RAM memory between locations 202:100 and 202:157 is as described in Example 2 above, i.e., 000 through 057 octal. With the motion control cable removed, position the tape in the cassette player to the beginning of the leader tone that was recorded in Example 2. Replace the motion control cable, adjust the volume setting and depress the play button.

The command register should be initialized just as it was in Example 3. The starting address register and the data block length should be set just as they were in Example 2. However, the secondary command register (register C) needs to be configured for a verify operation. Bit 6 is the verify bit, therefore, the C register should have the following pattern:

bit 7	bit 6	bit 5	bit 4	bit 3	bit 2	bit 1	bit 0
0	1	0	0	0	0	0	0

This pattern corresponds to the octal number 100 or the hex 40.

The results of verify operations are stored at location 203:365 octal or 83F5 hex and has the label DERR which stands for Data ERRor. The COPE software writes in this location only if it encounters data on the tape which is different from corresponding data in memory. If an error is found, COPE loads DERR with some non-zero value. Before a verify operation, DERR should be initialized to zero. At the end of a verify operation, DERR

should be checked. If DERR still has the value zero, then the verify was successful. If it is a non-zero, a compare error was detected and the data on the tape may be incorrect.

The following program contains the necessary instructions to verify the tape data against the data stored between locations 202:100 and 202:157 octal.

Octal			
202:054 202:057 202:061 202:063 202:066 202:071 202:074	061 370 202 076 202 016 100 021 060 000 041 100 202 315 012 201 303 074 202	LXI SP,202:370 MVI A,202 MVI C,100 LXI D,60 LXI H,202:100 CALL COPE JMP 202:074	INITIALIZE STACK POINTER INITIALIZE COMMAND REGISTER LOAD C WITH VERIFY COMMAND INITIALIZE BYTE COUNT REG INITIALIZE BLOCK STARTING ADDR EXECUTE COMMAND DYNAMIC HALT
Hex 822C 822F 8231 8233 8236 8239 823C	031 F8 82 3E 82 0E 40 11 30 00 21 40 82 CD OA 81 C3 3C 82	LXI SP,82F8 MVI A,82 MVI C,40 LXI D,30 LXI H,8240 CALL COPE JMP 823C	

Examine location 203:365 octal (83F5 hex). This is the Data Error location DERR. Store a zero in this location. Enter the above program starting at location 202:054 octal (822C hex). Start the program. The tape will start moving and continue until the compare operation is finished. When the tape stops, halt the computer and examine DERR. The value of this memory location should still be zero.

Reposition the tape to the beginning tone leader and then examine location 202:120. Change this value from 20 to 377. Restart the program at location 202:054 octal (822C hex). Again the tape will start moving and continue until the verify operation has been completed. When the tape stops, halt the computer and re-examine location 203:365 octal (83F5 hex). DERR should now have a value other than zero (357 in this instance).

#### BYTE ORIENTED DATA TRANSFERS

In the applications discussed previously, the data managed by COPE was always in blocks many bytes long and transferred as a whole rather than a single byte at a time. There are applications, however, where the length of a block of data on a tape is not known before hand. The end of a block of data is usually marked with a special code which corresponds to an "end of file" mark. In this situation, a block transfer is not practical. Data must be transferred a byte at a time and each byte must be tested to see if the transfer is finished.

In effecting this type of transfer, COPE will still be transferring blocks of data. However, each block of data will be only one byte long and the command register will be modified so that the tape motion is not turned off in between transfers.

For a read operation from, say, channel 1 where the tape is <u>not</u> to be turned off after the command has finished, the command register has the pattern:

bit 7	bit 6	bit 5	bit 4	bit 3	bit 2	bit 1	bit 0
0	0	0	0	0	0	1	0

As with a normal read, the C register is initialized to zero.

Assuming that the file in question is made up of ASCII characters which have values of octal 40 or greater, the ASCII code 004, call EOT (end of tape), will be used for an "end of file" mark for the data. The following program will transfer data from the tape cassette unit to memory starting at location 0.

#### Octal

203:000	061 370 203	START	LXI	SP,203:370	INITIALIZE STACK POINTER
203:003	041 000 000		LXI	Н,0	INITIALIZE DATA POINTER
203:006	345	LOOP	PUSH	H	SAVE THE DATA POINTER
203:007	021 001 000		LXI	D,1	SET UP FOR 1 BYTE TRANSFER
203:012	076 002		MVI	A,2	LOAD COMMAND REGISTER
203:014	016 000		MVI	C,0	INITIAL SECONDARY CMD REG
203:016	041 047 203		LXI	H,TEMP	LOAD BYTE TRANSFER ADDR
203:021	315 012 201		CALL	COPE	GET BYTE OF DATA FROM TAPE
203:024	072 047 203		LDA	TEMP	GET TAPE DATA BYTE IN ACC
203:027	376 004		CPI	EOT	COMPARE END OF FILE MARK
203:031	341		POP	H	RESTORE STACK/GET DATA PNTR
203:032	312 042 203		JZ	DONE	TEST FOR END OF FILE
203:035	167		MOV	M,A	MOVE DATA TO MEMORY
203:036	043		INX	H	INCREMENT DATA POINTER
203:037	303 006 203		JMP	L00P	GO GET MORE DATA
203:042	076 000	DONE	MV I	A,0	
203:044	323 004		OUT	TAPE	TURN OFF TAPE MOTION
203:046	311		RET		RETURN
203:047	000	TEMP	000		DATA CELL WHERE COPE
-					TRANSFERS DATA BYTE

Hex				
8300	31 F8 83	START	LXI	SP,83F8
8303	21 00 00		LXI	H,0
8306	E5	LOOP	PUSH	Н
8307	11 01 00		LXI	D, 1
830A	3E 02		MVI	A,2
830C	0E 00		MV I	C,0
830E	21 27 83		LXI	H, TEMP
8311	CD OA 81		CALL	COPE
8314	3A 27 83		LDA	TEMP
8317	FE 04		CPI	EOT
8319	E1		POP	Н
831A	CA 22 83		JZ	DONE
831D	77		MOV	M,A
831E	23		INX	H
831F	C3 06 83		JMP	LOOP
8322	3E 00	DONE	MV I	A,0
8324	D3 04		OUT	TAPE
8326	<b>C</b> 9		RET	
8327	00	TEMP	000	

#### THE CHECK SUM SOFTWARE

Oncedata is stored on some bulk magnetic media such as a disk or tape, it is always possible that through some fault in the media or from an accidental exposure to an intense magnetic field, the original data has been altered. It is also possible that the read electronics of the device(s) which transfer data from the media to the computer's memory are faulty or even worse periodically faulty. If a program stored on tape is read into memory with errors, it probably will not run, and so it is loaded again and usually the second or third time it goes in correctly. This is annoying but not disastrous. However, if a payroll file were to be read in with errors such that several employees were in effect given large raises, this might be considered disastrous. The field of error correcting and detecting codes and algorythms has received a good deal of study and experimentation for these and other reasons. Error correcting codes are quite complex and are beyond the scope of the ROM software on the I/O board. However, there is a short routine included in the ROM called CHECK whose starting address 201:115 octal (814D hex). This routine can be used to detect a wide variety of read errors which may occur in transferring data from the audio cassette player to the computer's memory. This section is devoted to explaining how to use the subroutine CHECK to enhance the integrity of data stored on audio cassette tapes.

CHECK is used to generate a "check sum" over a data block which is to be transferred to or from memory. This check sum is two bytes long. The most straightforward way to use CHECK is as follows:

(1) Just before writing a data block out on tape, generate a check-sum word over the data block with CHECK and then append these two bytes to the end of the data.

Octal

- (2) Write augmented data block on the tape.
- (3) Whenever the block of data is read into the computer, generate another check sum over all the data except the last two bytes.
- (4) Compare the new check sum with the original check sum which is stored in the last two bytes of the data block.

The CHECK subroutine generates check sums over blocks of data. Just as with the COPE software, CHECK requires that the starting address of the data block be in the H-L register pair and that the block length be in the D-E register pair. When CHECK is through, the check sum is in the H-L register pair -- high order byte in H and low order byte in L. The D-E register pair has the address of the next location above the data block. This makes it easy to append the check sum to a data block or to compare a new check sum with an old one. Below are two examples of subroutines which use CHECK. The first appends a check sum to a data block and the second compares a new check sum with an old one at the end of the block. The zero flag is set if the two check sums agree. In each case, the data block is 4K (4096) bytes long and starts at location 0.

01.1.000.000	400540	1 3/1		
041 000 000 021 000 020 315 115 201 353	APPEND:	XCHG	H,0 D,20:000 CHECK	INITIALIZE STARTING ADDR INITIALIZE BLOCK LENGTH COMPUTER CHECK SUM XCHG CHECK SUM W/ ADDR
163 043 162		MOV INX MOV	M,E H M,D	STORE LOW ORDER BYTE OF CS INCREMENT ADDR POINTER STORE HIGH ORDER BYTE OF CS
311		RET	п, о	RETURN
Hex				
21 00 00 11 00 10 CD 4D 81 EB	APPEND:	LXI	H,0 D,1000 CHECK	
73 23		MOV INX	M,E H	
72 C9		MOV RET	M, D	
Octa1			11	
041 000 000 021 000 020 315 115 201 353 173 276 300 043 172 276 311	COMPAR	LXI	H,0 D,20:000 CHECK A,E M H A,D	INITIALIZE STARTING ADDR INITIALIZE BLOCK LENGTH COMPUTE NEW CHECK SUM XCHG CHECK SUM W/ ADDR MOVE NEW LOW BYTE TO A COMPARE W/ OLD LOW BYTE RETURN IF ERROR INCREMENT ADDR MOVE NEW HIGH BYTE TO A COMPARE W/ OLD HIGH BYTE RETURN

# Principles of Operation

Hex			
21 00 00	COMPAR	LXI	Н,О
11 00 10		LXI	D,1000
CD 4D 81		CALL	CHECK
EB		XCHG	
7B		MOV	A,E
BE		CMP	М
CO		RNZ	
23		INX	Н
7A		MOV	A,D
BE		CMP	M
C9		RET	

#### THE SOFTWARE UART

#### INTRODUCTION

The serial port of the I/O board has been designed to interface to a teletype with a paper tape reader and any serial RS232 device. The ADM3-A terminal from Lear Seigler or the Hazeltine 1500 are examples of relatively inexpensive RS232 video terminals which interface directly to the serial port.

Many I/O boards have a serial port -- it is a very useful interface to have on an I/O board. However, most I/O boards with serial ports have an integrated circuit called a UART connected to the port. It takes care of the serial-to-parallel conversion for incoming data from a terminal and the parallel-to-serial conversion for outgoing data from the computer to the terminal. The Speakeasy does not have a UART; instead, it has several pieces of software in its PROM which together constitute a software UART. This intelligent parallel-to-serial and serial-to-parallel converter has a great deal of flexibility and in many ways is superior to most hardware UARTS. Unfortunately, it is not more flexible in all respects.

An integrated circuit UART is actually two separate pieces of hardware in a single package. The parallel-to-serial conversion and the serial-toparallel conversion functions of a UART are independent of each other. On many UARTS these two functions even have provision for separate clocks so that the two processes can take place at different rates. The parallel-toserial and serial-to-parallel functions of the software UART are also separate pieces of code. Unlike the hardware UART, however, only one of these processes can take place at a given time since the CPU itself is the prime mover in the process. The consequence of this is that for slow devices such as a teletype it is possible that during output, an input character from from the keyboard may be missed. This would be a serious problem if it were not for an important option that is built into the input (serial-to-parallel) software. The input software has the ability to echo input as it is collected so that during input it is not necessary to do any output (parallel-to-serial). At data rates of 300 baud and better, characters are collected (serial-to-parallel) and echoed (parallel-to-serial) faster than most humans can type so that it is virtually impossible to miss input during output. At these speeds the two separate processes of the software UART appear to someone at a terminal to be occuring simultaneously.

#### INPUT

The input software of the PROM UART allows for several options while it is collecting input data. A paper tape reader of a TTY can be turned on and characters can be echoed while they are being collected. The B register of the CPU is used to enable these options. The starting address of the serial input routine is 201:157 (816F hex) and is executed using a CALL INPUT instruction. The input character is returned in the D register. The input routine uses all the CPU registers and six levels of the stack.

#### **OUTPUT**

The output routine of the software UART expects that the parallel data to be sent out to the terminal is in the accumulator. The starting address of the output routine is 201:263 (81B3 hex). As with the input routine all the CPU registers are used along with four levels of the stack.

#### THE SPEED CONSTANT SCON

The software UART can run at any speed between 0 and 2400 baud (bits per second). The baud rate is set by initializing locations 203:363 and 203:364 (83F3 and 83F4). 203:363 has the label SCON which stands for Serial CONstant. Both SCON and SCON+1 must be initialized before a valid baud rate for the software UART is determined. The formula for calculating SCON and SCON+1 is given below:

$$SCON = \frac{1,000,000}{K \times (baud rate)}$$

For baud rates less than 600, K has the value 53 (decimal). For rates more than 600, K should be increased to 57. Below are some values for SCON for the more common baud rates:

Baud Rate	SCON+1,	SCO	1	
110	000:253	(00	AA	hex)
300	000:075	(00	3D	hex)
1200	000:016	(00	OE	hex)
1800	000:011	(00	09	hex)
2400	000:006	(00	06	hex)

Some terminals may require a slight change from the values listed above which are average quantities.

#### Examples

Octal

Listed below is a short program ideal for testing the input and output circuitry of the serial port of the I/O board. This is a very low level routine which does not use SCON and has the effect of making a half duplex terminal out of a full duplex one.

00	. Ca i	пе	:X		
202:002	333 005 323 005 303 000 202	8202	DB 05 D3 05 C3 00 82	START	 GET THE INPUT ECHO INPUT RT

Ц.,,

When the terminal is connected to the serial port and this program is running, whenever a key is pressed on the keyboard, it will be echoed on the screen or by the printer.

When bit 0 of register B is a zero, the input routine will automatically echo input. When bit 7 of register B is a 1, the TTY's paper tape reader will be turned on until a character is collected. It is always turned off when a character is collected even if some other piece of code has turned the reader on. Following is a program which exercises both the input and output routines. The speed constant is set for 300 baud.

Octal					
202:000 202:003 202:006	061 000 203 041 075 000 042 363 203	START	LXI LXI SHLD	SP,203:000 H,75 SCON	INITIALIZE STACK POINTER INITIALIZE SCON
202:011 202:013 202:016	006 002 315 157 201 172	L00P		B,1 INPUT	SUPPRESS AUTOMATIC ECHO GO GET INPUT
202:017 202:022	315 263 201 303 011 202		MOV CALL JMP	A,D SROUT LOOP	TRANSFER INPUT CHARAC TO A GO ECHO INPUT GO GET MORE INPUT
Hex					
8200 8203 8206	31 00 83 21 3D 00 22 F3 83	START	LXI LXI SHLD	SP,8300 H,3D SCON	
8209 820B 820E	06 01 CD 6F 81 7A	L00P	MVI CALL MOV	B,1	
821F 8212	CD B3 81 C3 09 82			SROUT LOOP	

When calling SROUT, the output routine, the character to be printed must be in the A register. When calling INPUT, the character entered from the keyboard is returned to the D register. The program below will perform just like the previous example. The difference between the two is that the CALL SROUT is missing but the INPUT routine is now enabled to echo characters as they are collected.

Octal				
202:000 202:003 202:006 202:011 202:013 202:016	061 000 203 041 075 000 042 363 203 006 000 315 157 201 303 011 202	START LOOP	LXI LXI SHLD MVI CALL JMP	SP,203:000 H,75 SCON B,0 INPUT LOOP
Hex				
8200 8203 8206 8209 820B 820E	31 00 83 21 3D 00 22 F3 83 06 00 CD 6F 81 C3 09 81	START LOOP	LXI LXI SHLD MVI CALL JMP	SP,8300 H,3D SCON B,0 INPUT LOOP

#### THE SERIAL STATUS FLIP-FLOP

Associated with the serial port is a status flip-flop which is accessed through the status port. This flip-flop is connected to bit 0 of input device 6. The serial status is set whenever a logical zero is present on either the TTY input or the RS232 input. Serial devices transmit a continuous stream of ones or rest bits when they are idle. When a serial device is ready to transmit data, it sends out one zero bit (start bit), seven or more data bits and one or more rest bits. The serial status flip-flop is a convenient signal to test whenever a program needs to know if a key (any key) has been struck on the keyboard. If a key has been struck,

333 006 IN 6 346 001 ANI 1

will reset the zero flag. If no key has been struck, the two instructions will set the zero flag. The status flip-flop is reset whenever an IN 5 instruction is executed.

#### THE PARALLEL PORT

The parallel port is perhaps the simplest of the four ports on the I/O board. It is a bidirectional port: data may be transmitted or received. There is a 1  $k\Omega$  pull-up resistor connected to pin #5 of J1 which disables the transmitters of the DM8835 transceivers. The default direction of the port is "input". This pin should be grounded if the port is to function as an output port. Associated with the output port is an 8-bit latch. This latch stores the most recent data sent to the parallel output port. There is a status bit of device 6 also associated with the parallel port. Pin #13 of J1 is connected to the clock input of a flip-flop. When a positive pulse is received at this pin, the status flip-flop for the port is set. When an IN 7 instruction is executed by the processor, the flip-flop is reset.

This port is called an inverting parallel port. The DM8835 transceivers buffering information from the outside world invert during both transmission and reception. Because of the higher noise margins, peripheral devices frequently send inverted data rather than positive data. When this is the case, the receivers of the DM8835s will automatically invert the data back to its natural positive state. If the peripheral transmits positive data, it will have the wrong polarity when an IN 7 instruction is executed. If this is the case, the IN 7 instruction should be followed by a CMA instruction to invert the data to its positive state. Below are two small routines —the first for input and the second for output — which illustrate the way a typical peripheral connected to the port should be accessed.

# Input Routine (Input data in A)

0ct	al	Hex								
202:100 102 104 107 111	333 006 346 002 312 100 202 333 007 311	8240 DB 06 8242 E6 02 8244 CA 40 82 8247 DB 07 8249 C9		IN 6 ANI 2 JZ PIN IN 7 RET	GET STATUS STRIP OFF PARA PORT BIT TEST IF DATA READY GET DATA/RESET STATUS RETURN					
	Output Routine (Output data in B)									
202:100 102 104 107 110 112 114	333 006 346 002 312 100 202 170 323 007 333 007 311	8240 DB 06 8242 E6 02 8244 CA 40 82 8247 78 8248 D3 07 824A DB 07 8246 C9		IN 6 ANI 2 JZ POUT MOV A,B OUT 7 IN 7 RET	GET STATUS STRIP OFF PARA PORT BIT TEST FOR OUTPUT READY MOVE OUTPUT DATA TO ACC TRANSMIT DATA RESET STATUS RETURN					

# Principles of Operation

A final example is the routine to access data from an Oliver OAE paper tape reader connected to the parallel port:

202:200	333 006	8280	DB 06	OPIN IN 6	GET STATUS
202:202	346 002	8282	E6 02	ANI 2	STRIP PARALLEL PORT BIT
202:204	312 200 202	8284	CA 80 82	JZ OPIN	TEST FOR INPUT READY
202:207	333 007	8287	DB 07	IN 7	GET INPUT DATA
202:211	057	8289	2 F	CMA	COMPLIMENT
202:212	311	828A	C9	RET	RETURN

#### BOOTSTRAP SOFTWARE AND FILE STRUCTURES

#### INTRODUCTION

Almost all of the modern micro-processor based computers use semi-conductor memory exclusively. This type of memory has many advantages but has one major short coming when compared with "magnetic core" memory --semi-conductor memory cannot remember the information stored in it when power was turned off. This means programs must be reloaded whenever power is turned off. The Speakeasy I/O board was designed to eliminate this problem.

There is a short routine contained in the COPE ROM called "BOOTS" which has the ability to automate the process of program loading and initialization. BOOTS stands for Bootstrap and the routine has been designed to perform system initialization by utilizing the COPE tape handling software. Specifically, this routine automatically loads a 256 byte record from the tape into the first page of the RAM on the I/O board. After this record is loaded, BOOTS forces an unconditional branch to the first location of this RAM. Just loading a short record and forcing a branch will not by itself perform system initialization. On the other hand, if the record that BOOTS loads is really a program which contains loading parameters pertaining to yet another program which follows immediately on the tape, then the process of initialization can be successfully completed.

The BOOTS routine is in ROM and always does the same task. The desired program that is to be loaded and started such as BASIC or ATE is on the second record of the tape and must be loaded. The bridge between BOOTS and the main program is the first 256 byte record/program on the tape that is loaded and then started by BOOTS. The structure of this record is the main topic of this section. This record has been designed both to facilitate program loading and to implement a file structure for records on audio cassette tapes.

The Speakeasy I/O board can be used in a variety of applications without any file structure. However, experience has shown that the board is even more effective when some type of file structure is implemented. Below is an example (suggestion) which can be ignored, followed, or used as a starting point for a more elaborate and effective file structure.

#### **HEADER STRUCTURE**

The 256 byte record that B00TS loads and that is in the front of any program on tape will be called a "header" or "record header." This header contains instructions for reading and writing records using COPE. It also contains information pertaining to the record that it proceeds. The information implements the file structure and will be discussed first. The instructions along with a "map" of the header will be presented afterwards.

The header contains the following file dependent information:

- (1) A checksum (2 bytes) for error detection.
- (2) The load address in memory (2 bytes) for the record.
- (3) The length of the record (2 bytes).
- (4) The record type (1 byte).
- (5) The tape cassette write channel number (1 byte).
- (6) The title type (1 byte).
- (7) The title (130 bytes).

Checksum has the label CHKSM+1 and contains the two byte checksum generated by CHECK when the record is originally created.

Load address has the label LODAD. If the record is a program, the starting address should agree with the load address.

Record length has the label LNGTH.

Record type has the label TYPE. Allowance has been made for up to 250 different record types. If TYPE contains an ASCII B, the record contains binary data or a program. If TYPE contains an ASCII S, the record consists entirely of ASCII characters and generally would be source code for a program.

Tape cassette write channel number has the label WUNIT. WUNIT contains 203, 205, or 211 (83, 85, 89 hex) and is in actuality a command to COPE to write on channel 1, 2, or 3. This byte is present to facilitate the reproduction (with perhaps minor changes) of records without the aid of any higher level software.

Title type has the label TITYP. Some records have no title or have an encoded title. If TITYP contains an ASCII B, the record has a binary title which consists of the first 130 bytes of the record in binary. If TITYP contains an ASCII S, the record has a title consisting of ASCII characters.

Title has the label TITLE. If the title type is an ASCII S, the title consists entirely of ASCII characters. It may be as long as 130 bytes. The last character of the title is an ASCII carriage return which is the octal number 15 (hex  $\emptyset D$ ).

Below are the instructions of the header program/record which uses the file structure information presented earlier to load and reproduce records.

# Octal

202:000	333 377	HEADR	IN SSW	GET SENSE SWITCHES
202:002	017		RRC	TEST FOR SSW Ø
202:003	332 000 203		JC HEADR	WAIT IF SSW Ø ON
202:006	061 066 203		LXI SP, TSTAK	INITIALIZE THE STACK
202:011	315 031 203		CALL LÓADR	CALL LOAD RECORD ROUTINE
202:014	302 014 203	STOP1	JNZ STOP1	TEST FOR ERRORS
202:017	333 377	NOGO	IN SSW	GET SENSE SWITCHES
202:017	007	NOGO	RLC	TEST FOR SSW 7
	-			
202:022	332 017 203		JC NOGO	WAIT IF SSW 7 ON
202:025	052 000 203		LHLD VERAD	GET FIRST ADDRESS OF RECORD
202:030	351		PCHL	BRANCH THERE
202:031	052 167 203	LOADR		GET THE LOAD ADDRESS
202:034	042 000 203		SHLD VERAD	SAVE FOR LATER VERIFY
202:037	315 062 203		CALL LREGS	LOAD UP THE REGISTERS
202:042	315 136 203		CALL COPER	GET READY TO READ RECORD
202:045	300		RNZ	RETURN IF STATUS ERROR
202:046	315 115 201		CALL CHECK	COMPUTE CHECK SUM
202:051	021 000 000	CHKSM		LOAD UP OLD CHECK SUM
202:054	174		MOV A,H	GET LOW BYTE OF CHECK SUM
202:055	272		CMP D	COMPARE WITH OLD LOW BYTE
202:056	300		RNZ	RETURN IF NO COMPARE
202:057	175		MOV A,L	GET HIGH BYTE OF CHECK SUM
	273		CMP E	COMPARE WITH OLD HIGH BYTE
202:060				
202:061	311		RET	RETURN
202:062	052 171 203	LREGS	LHLD LNGTH	LOAD THE D-E REGISTER
202:065	352		XCHG	PAIR WITH RECORD LENGTH
202:066	052 000 203		LHLD VERAD	LOAD H-L WITH STARTING ADDR
202:071	076 202		MVI A,202Q	INITIALIZE A & C WITH
202:073	016 000		MVI C,Ø	A READ COMMAND
202:075	311		RET	RETURN
202:076	315 104 202	TAPEW		GO WRITE HEADER & RECORD
202:101	303 101 202	STOP2	JMP STOP2	DYNAMIC HALT
202:104	315 062 202	TAPE	CALL LREGS	LOAD UP CPU REG FOR COPE
202:107	315 115 201		CALL CHECK	COMPUTE CHECK SUM
202:112	042 052 202		SHLD CHKSM+1	SAVE CHECK SUM
202:115	041 000 202		LXI H, HEADR	STARTING ADDR FOR HEADER
202:120	021 000 001		LXI D,1:000	LENGTH OF HEADER
202:123	072 174 202		LDA WUNIT	WRITE COMMAND
202:126	365		PUSH PSW	SAVE FLAGS & ACCUMULATOR
202:127	315 136 202		CALL COPER	GO WRITE THE HEADER
202:127	315 062 202		CALL LREGS	LOAD UP CPU REG FOR COPE
_			POP PSW	RESTORE FLAGS & ACC
202:135	361 265	COPER	PUSH PSW	SAVE FLAGS & ACC
202:136	365	CUPER	**	CLEAR ACC
202:137	257		XRA A	
202:140	062 365 203		STA DERR	INITIALIZE DATA ERROR CELL
202:143	062 366 203		STA SERR	INITIALIZE DATA ERROR CELL
202:146	361		POP PSW	RESTORE FLAGS & ACC
202:147	345		PUSH H	SAVE STARTING ADDRESS
202:150	335		PUSH D	SAVE BLOCK LENGTH
202:151	315 012 201		CALL COPE	GO DO THE TAPE COMMAND
202:154	331		POP D	RESTORE BLOCK LENGTH
202:155	341		POP H	RESTORE STARTING ADDRESS

## Bootstrap Software and File Structures

202:156	072 366 203	LDA SERR	GET THE STATUS ERROR CELL
202:161	107	MOV B,A	SAVE IN B
202:162	072 365 203	LDA DERR	GET THE DATA ERROR CELL
202:165	250	ORA B	MERGE ERROR DATA
202:166	311	RET	RETURN
202:167	000 000	LODAD DW O	LOAD ADDRESS WORD
202:171	000 000	LENGTH DW O	BLOCK LENGTH WORD
202:173	000	TYPE DW O	FILE TYPE
202:174	203	WUNIT DB 203Q	UNIT NUMBER & WRITE CMD
202:175	102	TITYP DB 102Q	TITLE TYPE
202:176	000	TITLE DS 202Q	TITLE AREA
Hex			

1		
п	ex	

8200	DB FF	HEADR	IN SSW
8202	0F		RRC
8203	DA 00 82		JC HEADR
8206	31 36 83		LXI SP, TSTAK
8209	CD 19 82		CALL LOADR
820C	C2 OC 82	STOP1	JNZ STOP1
820F	DB FF	NOGO	IN SSW
8211	07		RLC
8212	DA OF 83		JC NOGO
8215	2A 00 83		LHLD VERAD
8218	E9		PCHL
8219	2A 77 82	LOADR	LHLD LODAD
821C	22 00 82		SHLD VERAD
821F	CD 32 82		CALL LREGS
8222	CD 5E 82		CALL COPER
8225	CO		RNZ
8226	CD 4D 81		CALL CHECK
8229	11 00 00	CHKSM	LXI D,0
822C	7C		MOV A,H
822 D	ВА		CMP D
822E	CO		RNZ
822F	7D		MOV A,L
8230	ВВ		CMP E
8231	<b>C9</b>		RET
8232	2A 79 82		LHLD LNGTH
8235	EB		XCHG
8236	2A 00 83		LHLD VERAD
8239	3E 82		MVI A,82
823B	0E 00		MVI C,O
823D	<b>C9</b>		RET
823E	CD 44 82	TAPEW	CALL TAPE
8241	C3 41 82	STOP2	JMP STOP2
8244	CD 32 82	TAPE	CALL LREGS
8247	CD 4D 81		CALL CHECK
824A	22 2A 82		SHLD CHKSM

824D 8250 8253 8256	21 00 82 11 00 01 3A 7C 82 F5		LXI H, HEADR LXI D, 100 LDA WUNIT
8257	CD 5E 82		PUSH PSW CALL COPER
825A	CD 32 82		CALL LREGS
825D	F1		POP PSW
825E	F5	COPER	PUSH PSW
825F	AF		XRA
8260	32 F5 83		STA DERR
8263	32 F6 83		STA SERR
8266	F1		POP PSW
8267	E5		PUSH H
8268	D5		PUSH D
8269	CD OA 81		CALL COPE
826C	D1		POP D
826D	E1		POP H
826E	3A F6 83		LDA SERR
8271	47		MOV B,A
8272	3A F5 83		LDA DÉRR
8275	B0		ORA B
8276	<b>c</b> 9		RET

#### THE HEADER PROGRAM

Once the header program is in memory, it should be preserved for later use with other programs. Below are the steps to save a copy of the header program so that it will auto-load and not try to load any record that follows it. This copy should be used only when it is necessary to create a file structure for a program or a data record.

- (1) Enter the header program starting at location 202:000 (8200 hex).
- (2) Enter 101 (41 hex) at location 202:167.
- (3) Enter 202 (82 hex) at location 202:170.
- (4) Enter 001 (01 hex) at location 202:171.
- (5) Put a tape in the recorder connected to channel one and press record.
- (6) Start the header program at location 202:076 (823E hex).

The recorder will start and will write out the header record itself on tape. Next, the program will write exactly one byte from location 202:101 of the header program. When BOOTS is used to read in the header record, regardless of the position of the sense switches, the header program will load and then come to a dynamic halt. If none of the sense switches are on, this halt will occur at location 202:101.

#### AUTO-LOADING PROGRAMS

With the header program in memory, any other program can be preserved on tape with auto loading capability by following the steps below:

- (1) Enter the lowest address of the program at locations 202:167 and 202:170; the low order byte of the address is loaded at 202:167 (8277 hex) and the high order byte is entered at 202:170 (8278 hex).
- (2) Enter the length of the program at locations 202:171 and 202:172. The length is a two byte value to allow programs of any length to be preserved. The low order byte is stored at location 202:171 (8279 hex) and the high order byte at location 202:172 (827A hex).
- (3) Start the header program at location 202:076 (823E hex).

The header program will first record the header on the tape after a leader tone of about five seconds has been recorded. Another five second leader tone is recorded, followed by the main program.

An auto-loading program consists of two records: the first is the header record and the second is the program itself. The header record contains file parameter information and code to load the program which follows it. To load an auto-loading program, the bootstrap program BOOTS is executed. There are two sense switches used by the header program which allows a considerable degree of flexibility: sense switch 0 and sense switch 7.

Sense switch 0: If sense switch 0 is turned on while the header record is loading, the header program will not load the main record until sense switch 0 is turned off. This feature is very useful when a record needs to be placed in a different area than actually called for. After the header program loads and is waiting for sense switch 0 to be turned off, stop the program. Change locations 202:167 and 202:170 (8277 and 8278 hex) to the new load address. Restart the header program at location 202:000 (8200 hex). Turn on sense switch 7 just after restarting the header program.

Sense switch 7: When this switch is turned on, the header program will load the main record but will <u>not</u> allow the main program to execute. This is useful if it is necessary to set certain parameters of the main program before execution. It is also a necessary procedure if the main program's lowest address is not its starting address or if the main program is not really a program but a data file.

#### AN "INTEL HEX" LOADER

There is a substantial body of software available to the small computer user which is recorded on KC Standard cassette tapes and has a data format which conforms to the so-called Intellec Hex Paper Tape Format or "Intel Hex" format. The details of this format along with a program which will load "Intel Hex" formatted KC Standard tapes is presented below.

#### THE INTELLEC HEX PAPER TAPE FORMAT

In this format, <u>two</u> ASCII hexadecimal characters are used to represent a byte of data. Preceding the first data field and following the last data field, there must be a leader/trailer of at least 200 rest bits (approximately eight seconds of 2400 hz tone). Each record consists of six fields. The fields must occur in the order presented below.

1. RECORD MARK FIELD: Frame 0

The ASCII code for a colon (:) is used to signal the start of a record.

2. RECORD LENGTH FIELD: Frames 1 and 2

The number of data bytes in the record is represented by two ASCII hexadecimal digits in this field. The high order digit is in Frame 1. The maximum number of data bytes in a record is 255 (FF in hex). An end-of-file record contains two ASCII zeros in this field.

LOAD ADDRESS FIELD: Frames 3-6

The four ASCII hexadecimal digits in Frames 3-6 give the address at which the data is loaded. The high-order digit is in Frame 3, the low-order digit in Frame 6. The first data byte is stored in the location indicated by the load address; successive bytes are stored in successive memory locations. This field in an end-of-file record contains zeros or the starting address of the program.

4. RECORD TYPE FIELD: Frames 7 and 8

The two ASCII hexadecimal digits in this field specify the record type. The high order digit is in Frame 7. All data records are type 0; end-of-file records are type 1.

5. DATA FIELD: Frames 9 to 9 + 2 (record length) -1

A data byte is represented by two frames containing the ASCII characters 0-9 or A-F, which represent a hexadecimal value between 0 and FF (0 and 255 decimal). The high order digit is the first frame of each pair. There are no data bytes in an end-of-file record.

# 6. CHECKSUM FIELD: Frames 9 + 2 (record length) to 9 + 2 (record length) + 1

The checksum field contains the ASCII hexadecimal representation of the two's complement of the 8-bit sum of the 8-bit bytes that result from converting each pair of ASCII hexadecimal digits to one byte of binary, from the record length field to and including the last byte of the data field. Therefore, the sum of all the ASCII pairs in a record after converting to binary from the record length field to and including the checksum field is zero.

#### Format Example:

:130160000E0C119231CD40000E09119031CD40006E

:0A3170007E3196310100000092311B

:10317C0092310100963180008C31923100009631F1

:04318E0092319231B7

#### The Loader Program

This program will load Intel Hex format KC Standard tapes. If the loader encounters an error, it will stop at ERROR. Otherwise, the loader stops at DONE. This program must reside in RAM.

Octal				
203:000 203:003 203:006 203:010	061 200 203 315 020 203 076 000 323 004	START	LXI SP,TSTAK CALL READ MVI A,0 OUT 4	INITIALIZE STACK POINTER READ THE RECORD STOP THE CASSETTE PLAYER
203:012 203:015	302 012 203 303 015 203	ERROR DONE	JNZ ERROR JMP DONE	STOP HERE ON ERROR STOP HERE IF NO ERRORS
203:020 203:023	315 135 203 376 072	READ	CALL RBYTE CPI ":"	READ A BYTE FROM TAPE COMPARE WITH START OF REC
203:025 203:030	302 020 203 036 000		JNZ READ MVI E,0	WAIT FOR BEGINNING OF RECINITIALIZE CHECKSUM
203:032 203:035	315 075 203 120		CALL CHTR MOV D,B	GET A CHARACTER PAIR RECORD LENGTH TO D
203:036 203:041 203:042	315 075 203 140 315 075 203		CALL CHTR MOV H,B CALL CHTR	GET A CHARACTER PAIR HIGH ORDER ADDR BYTE TO H
203:045 203:046	150 315 075 203		MOV L,B CALL CHTR	GET ANOTHER CHARACTER PAIR LOW ORDER ADDR BYTE TO L
203:051 203:053	076 001 270		MVI A,1	GET RECORD TYPE FIELD COMPARE RECORD TYPE TO END OF RECORD
203:054 203:055	310 315 075 203	L00P	RZ CALL CHTR	IDENTIFIER GET A DATA BYTE
203:060 203:061	167 043		MOV M,A INX H	STORE IN MEMORY ADVANCE ADDRESS POINTER
203:062	025		DCR D	DECREASE LENGTH COUNTER

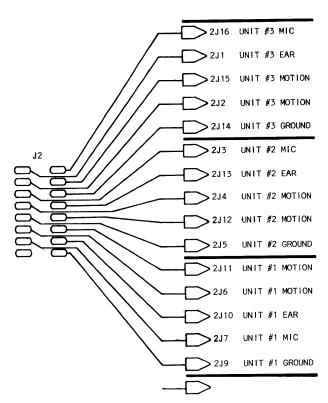
203:063 203:066	302 055 203 315 075 203			LOOP CHTR	GET MORE DATA IF NON-ZERO GET CHECKSUM
203:071 203:072	300 303 020 203		RNZ JMP	READ	STOP IF CHECKSUM ERROR READ NEXT RECORD IN
203:075 203:100 203:103 203:104 203:105 203:106 203:107 203:110 203:113 203:116 203:117 203:120 203:121	315 135 203 315 123 203 027 027 027 027 107 315 135 203 315 123 203 260 107 303 137 311	CHTR	CALL RAL RAL RAL MOV	B,A RBYTE	READ A BYTE FROM TAPE CONVERT TO HEX MULTIPLY BY HEXADECIMAL TEN SAVE IN B GET LOW ORDER ASCII DIGIT CONVERT TO HEX MERGE WITH HIGH ORDER BYTE SAVE IN B UPDATE CHECKSUM SAVE NEW CHECKSUM RETURN
203:123 203:125 203:127 203:131 203:132 203:134 203:135	076 000 326 060 376 012 330 326 007 311 305	HEX RBYTE	MVI SUI CPI RC SUI RET PUSH		GET THE DATA BYTE SUBTRACT BASE COMPARE WITH TEN RETURN IF LESS THAN TEN SUBTRACT OFFSET RETURN SAVE THE
203:136 203:137 203:140 203:143 203:146 203:150 203:151 203:154 203:155 203:156 203:157	325 345 041 124 203 021 001 000 076 002 113 315 012 201 341 321 301 311		PUSH PUSH LXI LXI MVI MOV CALL POP POP POP RET		STATE OF THE LOADER INITIALIZE STORAGE POINTER INITIALIZE BYTE COUNT INITIALIZE CMD REG INITIALIZE C TO ZERO READ A BYTE FROM TAPE RESTORE THE STATE OF THE LOADER RETURN
Hex					
8300 8303 8306 8308 830A 830D	31 80 83 CD 10 83 3E 00 D3 04 C2 0A 83 C3 0D 83	START . ERROR DONE	LXI CALL MVI OUT JNZ JMP	SP,TSTAK READ A,O 4 ERROR DONE	
8310 8313 8315 8318 831A 831D 831E 8321	CD 5D 83 FE 3A C2 10 83 1E 00 CD 3D 83 50 CD 3D 83 60	READ	CPI JNZ MVI CALL MOV	RBYTE ':'' READ E,O CHTR D,B CHTR H,B	

8322 8325 8326 8329 832B 832C	CD 3D 68 CD 3D 3E 01 B8 C8			CALL MOV CALL MV I CMP RZ	CHTR L,B CHTR A,1 B
832D 8330 8331 8332 8333	CD 3D 77 23 1D C2 2C		LOOP	CALL MOV INX DCR JNZ	CHTR M,A H D LOOP
8336 8339 833A	CD 3D CO	83		CALL RNZ	CHTR
833D 8340 8343 8344 8345 8346	CD 5D CD 53 17 17 17 17	83	CHTR	JMP CALL CALL RAL RAL RAL RAL	HEX
8347 8348 834B 834E 834F 8350 8351 8352	47 CD 5D CD 53 BO 47 83 5F C9	83 83		MOV CALL CALL ORA MOV ADD MOV RET	B,A RBYTE HEX B B,A E
8353 8355 8357 8359 835A	3E 00 D6 30 FE 0A D8 D6 07		HEX	MVI SUI CPI RC SUI	A,0 30H 10
835C 835D 835E 835F 8360 8363 8366 8366 8366 836D 836E 836F	11 01 3E 02 4B	83 00 81	RBYTE	RET PUSH PUSH LXI LXI MOV CALL POP POP RET	B D H H,HEX+1 D,1 A,2 C,E COPE H D B

#### INTERFACE CABLES AND CONNECTION PROTOCOL

#### THE TAPE CASSETTE PORT

J2 is the cassette interface connector. It is a 16-pin DIP (dual-in-line) socket. Fifteen of the sixteen pins are used. Five connections are provided for each tape channel. The pin-out for this connector is detailed below.

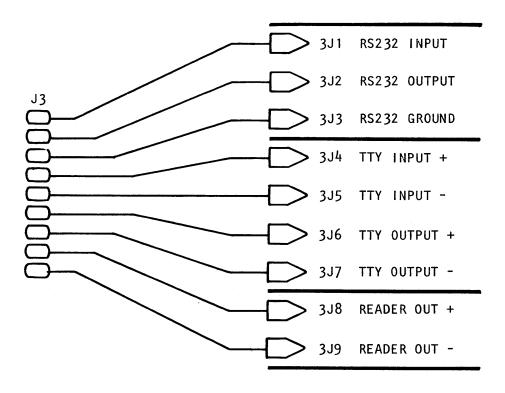


The input signal should come from the external speaker output of the cassette player/recorder. The output signal carries a wave form of approximately four volts peak and is much too high for the microphone input of the recorder. Instead, it should go to the auxiliary input. The two motion control outputs should be connected to the remote input of the cassette player/recorder. There are two choices for cables: flat ribbon or co-axial. If flat cable is used, it should terminate at one end with a DIP header (3M part #3416-0000 or equivalent). If co-axial cables are used, a 16-pin component carrier is a convenient way to terminate. There is one common ground connection which should be connected to both the auxiliary input ground connection and the external speaker output ground connection. (A 16-conductor flat cable with a DIP header is available from Morrow's to facilitate connecting J2 to a cassette player/recorder.)

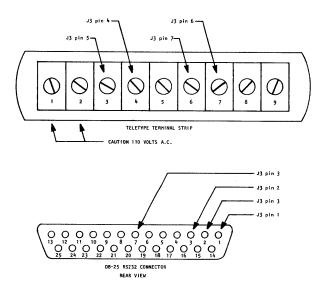
There are cassette players which have a ground loop between the ground connection of the external speaker jack and the auxiliary input jack. A ground loop between two common ground connections exists when there is a significant impedence between the two points. In a cassette player, ground loops generally are caused by component layout anomalies. Ground loop problems manifest themselves by hum or interference when both the auxiliary plug and the speaker plug are plugged in at the same time. If there is a noticeable background hum on tape data files, it may be caused by a ground loop in a cassette player. The problem can be averted by plugging in only one of the two signal plugs at a time: speaker plug during tape reads or auxiliary input during tape writes.

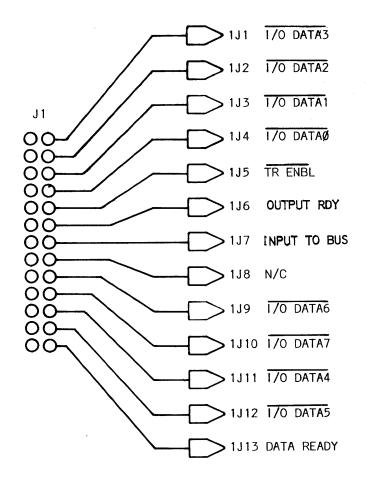
## THE SERIAL PORT

J3 is the serial interface connector. It is a 16-pin DIP socket. Nine of the sixteen pins are used. Three connections are dedicated to the RS232 section of the serial port and six are devoted to the TTY section of the port. The pin-out for this connector is detailed below.



If a TTY is connected to the port, the RS232 inputs should be left open. On the other hand, if an RS232 device is connected to the port, the TTY inputs must be connected. That is, pins 3 and 4 of J3 must be connected together when an RS232 device is connected to pins 1, 2 and 15 of J3. The two layouts below detail how the connectors of a TTY or an RS232 device (DB25) should be wired to J3. A 16-conductor flat cable is available from Morrow's to facilitate connecting J3 to a TTY or RS232 device.





## THE PARALLEL PORT

J1 is the parallel port connector. It is a 26-pin array designed to mate with a flat cable terminated with a 26-connection socket (3M part # 3399-0000 socket or equivalent). The pin out for this connector is detailed below. A 26-conductor cable with a 3M socket is available from Morrow's to facilitate connecting the parallel port to an I/O device.

	b1 J1 pin 1	1/0 DATA Ø J1 pin 4
	b2 J1 pin 2	1/0 DATA 1 J1 pin 3
	b3 J1 pin 3	1/0 DATA 2 J1 pin 2
	b4 J1 pin 4	1/0 DATA 3 J1 pin 1
	b5 J1 pin 5	1/0 DATA 4 J1 pin 11
	b6 J1 pin 6	I/O DATA 5 J1 pin 12
	b7 J1 pin 7	1/0 DATA 6 J1 pin 9
	b8 J1 pin 8	1/0 DATA 7 J1 pin 10
		DATA READY J1 pin 13
CHERRY "PRO" KEYBOARD INT	ERFACE EXAMPLE	
	0 <b>∮</b> 1/0 pin 1	I/O DATA Ø J1 pin 4
	D1 1/0 pin 16	1/0 DATA 1 J1 pin 3
	02 1/0 pin 2 .	1/0 DATA 2 J1 pin 2
	03 1/0 pin 15	1/0 DATA 3 J1 pin 1
	D4 1/0 pin 3	1/0 DATA 4 J1 pin 11
	D5 1/0 pin 14	1/0 DATA 5 J1 pin 12
	D6 1/0 pin 4	1/0 DATA 6 J1 pin 9
	D7 1/0 pin 13	1/0 DATA 7 J1 pin 10
	RDA 1/0 pin 7	DATA READY J1 pin 13
	ACK 1/0 pin 5	INPUT TO BUS J1 pin 7
OLIVER PAPER TAPE READER INTER	RFACE EXAMPLE	
	PARALLEL PORT IN SPEAKEASY I/O BO	TERFACE EXAMPLES A THINKER TOY PRODUCT FROM MORROW'S

#### PARTS LIST

```
8" x 10" glossy photograph
1
       5" x 10" printed circuit board
1
       24\Omega ¼ watt resistors
3
                                           (red-yellow-black)
       240\Omega \frac{1}{2} watt resistors
                                           (red-yellow-brown)
3
2
       470\Omega 1 watt resistors
                                           (yellow-purple-brown)
9
       1k\Omega \pm watt resistors
                                           (brown-black-red)
                                           (brown-green-red)
 1
       1.5k\Omega
       3.3k\Omega
                    11
                                           (orange-orange-red)
 3
5
       4.7k\Omega
                                            (yellow-purple-red)
       10k\Omega
                    11
2
                                            (brown-black-orange)
                                            (brown-green-orange)
 3
       15k\Omega
       22k\Omega
 1
                                            (red-red-orange)
2
       27k\Omega
                                            (red-purple-orange)
 4
       47k\Omega
                                            (yellow-purple-orange)
 1
        330kΩ
                                            (orange-orange-yellow)
        .001 µfd disk capacitor
 1
 4
        .01
        .1 μfd mylar capacitor
 1
2
        .82 μfd tantulum capacitors
 3
       39 µfd
17
       by-pass capacitors*
        1N914/4820-0201 signal diodes
7
 2
        1N4742 12 v. 1 watt zener diodes
        1N4585 rectifying diode
 1
 3
        Gordos 761A52 relays
```

<sup>\*</sup>by-pass capacitor values may vary from .01  $\mu fd$  to .1  $\mu fd$ 

- 3 2N3906 PNP transistors
- 1 2N3904 NPN transistor
- 17 low profile 16-pin sockets
- 6 low profile 14-pin sockets
- 1 low profile 20-pin socket
- 1 26 conductor right angle male header (position J1)
- 1 16 conductor right angle male header with shell & female connectors (J2)
- 7 or 9 conductor right angle male header with matching shell and female connectors (J3)
- 1 heat sink
- 1 set machine nut & screw
- 1 74LS00 quad 2-input NAND gate
- 1 74LS04 hex inverter
- 1 7406 hex open collector inverter/buffer
- 1 74LS37 quad 2-input NAND buffer
- 2 74LS74 dual D-type flip-flop
- 1 74LS139 dual 1 of 4 decoder
- 1 74LS155 dual 1 of 4 decoder
- 1 74LS175 quad dual rail register
- 1 74LS273 octal latch
- 2 74LS367 2-4 tri-state\* buffer with enables
- 4 2112/5039927A 4x256 N-MOS RAM
- 2 MMI 6306 4x512 PROM (Monolithic Memories)
- 1 MMI 6301/82S129 4x256 PROM
- 2 DS8835 inverting tri-state\* transceivers
- 2 8T97 2-4 tri-state\* buffer/driver with enables
- 1 741 operational amplifier
- 1 LM311 comparator

<sup>\*</sup>tri-state is a trade mark of National Semiconductor

## ASSEMBLY INSTRUCTIONS

DO NOT INSTALL OR SOLDER ANY PARTS UNTIL YOU HAVE READ THESE INSTRUCTIONS SEVERAL TIMES AND HAVE FULLY DIGESTED THE INFORMATION!

CAUTION -- DO NOT SOLDER OR CLIP COMPONENT LEADS WITHOUT USING SAFETY GLASSES!

## INSPECTION

Use the Parts List to make sure that there are no missing items in your kit. Please notify us of any shortages. Be sure to check for missing parts <u>before</u> you start to assemble.

## COMPONENT LEAD WIDTHS

Bend the leads with the plastic bending block in your kit. Be sure that the leads are bent to the proper width <u>before</u> the part is inserted. Properly bent components solder easier and give the finished kit a professional appearance.

- 1. All the  $\frac{1}{4}$  watt resistors have lead widths of .5 inches.
- 2. The  $\frac{1}{2}$  watt resistors R38 and R39 have lead widths of .5 inches.
- 3. The  $\frac{1}{2}$  watt resistor at R28 has a lead width of .6 inches.
- 4. The two 1 watt resistors have lead widths of .95 inches.
- 5. The two .82  $\mu$ fd tantulum capacitors C5 and C8 have lead widths of  $\frac{1}{2}$  inch.
- 6. The three 39  $\mu$ fd tantulum capacitors C9, C10, and C11 have lead widths of .6 inches.
- 7. The CR11 diode has a lead width of .7 inches; all other diodes have lead widths of  $\frac{1}{2}$  inch.
- 8. The transistors and ceramic disk capacitors have partially preformed leads and may be installed without preparation.

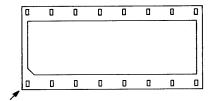
## SOCKETS

A socket is furnished for every integrated circuit. It is important that you use the sockets; otherwise, a defective part will be extremely difficult to replace.

NO REPAIR OR SERVICE WILL BE PERFORMED ON A KIT WHICH HAS HAD INTEGRATED CIRCUITS SOLDERED TO THE CIRCUIT BOARD.

## PARTS ORIENTATION

In all references throughout the instructions, the convention used is that the gold edge connector is the <u>bottom</u> of the board. Orientation identification is molded into the plastic of the sockets and is illustrated below:



This orientation mark identifies where pin #1 of the integrated circuit is to be positioned when inserted into the socket. The sockets should be inserted in the board so that the orientation mark is in the <u>lower left</u> hand corner.

Orientation of the transistors, tantulum capacitors, diodes and voltage regulator is specified in the component layout drawing. It is advisable to study this drawing and the  $8 \times 10$  glossy photograph carefully before building the kit. Refer to both during parts installation.

## SOLDERING AND SOLDER IRONS

The most desirable soldering iron for complex electronic kits is a constant temperature soldering iron with an element regulated at 650° F. The tip should be fine so that it can be brought in intimate contact with the pads of the circuit board. Both Unger and Weller have excellent products which fit the above requirements.

There are three important soldering requirements for building this kit:

- 1. Do not use an iron that is too cold (less than  $600^{\circ}$  F) or too hot (more than  $750^{\circ}$ ).
  - 2. Do not apply the iron to a pad for extended periods.
  - 3. Do not apply excessive amounts of solder.

The proper procedure for soldering components to the circuit board is as follows:

- 1. Bring the iron in contact with both the component lead and the pad.
- 2. Apply a <u>small</u> amount of solder at the point where the iron, component lead, and pad all make contact.
- 3. After the initial application of solder has been accomplished with the solder flowing to the pad and component lead, the heat of the iron will have transferred to both the pad and the lead. Apply a small amount of additional solder to cover the joint between the pad and the lead. DO NOT PILE SOLDER ON THE JOINT! EXCESSIVE HEAT AND SOLDER CAUSE PADS AND LEADS TO LIFT FROM THE CIRCUIT BOARD. EXCESSIVE SOLDER IS THE PRIMARY CAUSE FOR BOARD SHORTS AND BRIDGED CONNECTIONS.

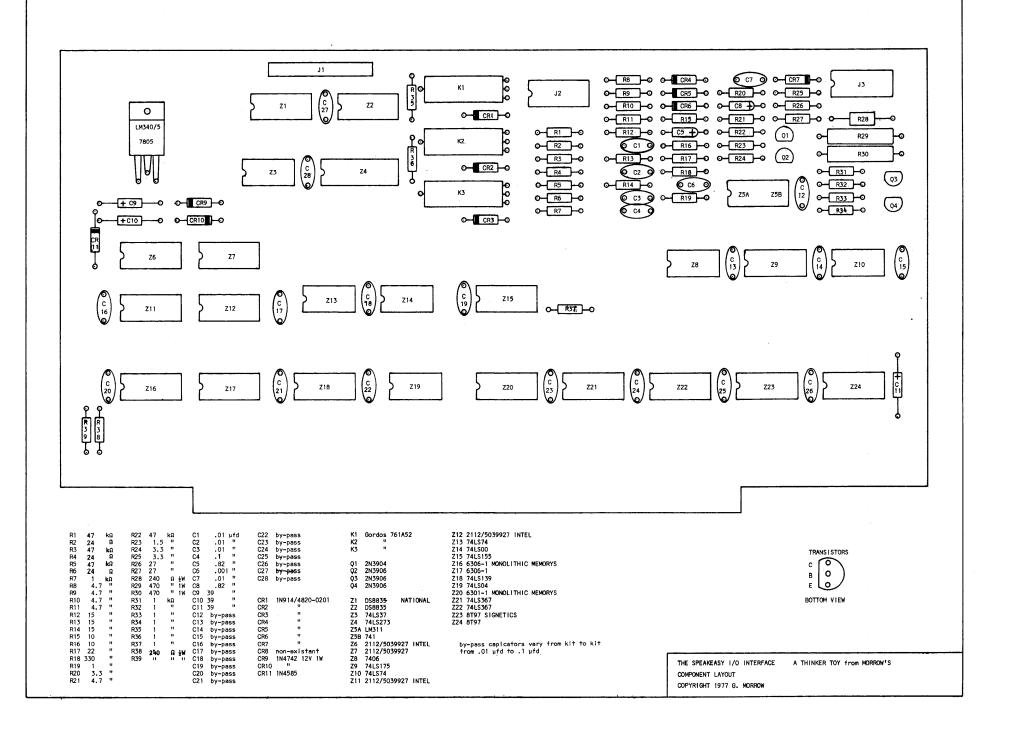
## PARTS INSTALLATION

Before installing parts, bend the leads of the resistors, diodes and tantulum capacitors to the proper lengths. After a series of parts have been installed in the board, bend the leads slightly to hold them in place, solder the leads and trim the excess lead lengths before proceeding to the next step.

## Install:

 R38-R39	$240\Omega$ ½ watt (bottom left)	(red-yellow-brown)
 CR9-CR10	1N4742 12 v. 1 watt zener	Check for orientation!
 R35-R36	1k $\Omega$ $\frac{1}{4}$ watt (top)	(brown-black-red)
 CR1-CR3	1N914/4820-0201	Check for orientation!
 R1, R3, R5	$47k_{\Omega}$ $\frac{1}{4}$ watt	(yellow-purple-orange)
 R2, R4, R6	$24_\Omega$ $\frac{1}{4}$ watt	(red-yellow-black)
 R7, R37	1kΩ ½ watt	(brown-black-red)
 R8-R11	4.7k $_{\Omega}$ $^{1}$ watt	(yellow-purple-red)
 R12-R14	15k $\Omega$ $\frac{1}{4}$ watt	(brown-green-orange)
 CR4-CR6	1N914/4820-0201	Check for orientation!
 R15-R16	$10k_\Omega$ $^{\frac{1}{4}}$ watt	(brown-black-orange)
 C5	.82 $\mu$ fd tantulum	Check for orientation!
 R17	$22k_\Omega$ $^{1\over 4}$ watt	(red-red-orange)
 R18	$330$ k $\Omega$ $\frac{1}{4}$ watt	(orange-orange-yellow)
 R19	1kΩ ¼ watt	(brown-black-red)
 R20, R24, R25	$3.3$ k $_\Omega$ $^{1\over 4}$ watt	(orange-orange-red)
 C8	.82 $\mu$ fd tantulum	Check for orientation!
 R2 1	4.7k $\Omega$ ½ watt	(yellow-purple-red)
 R22	$47k_{\Omega}$ $^{1\over 4}$ watt	(yellow-purple-orange)
 R23	1.5k <sub>Ω</sub> ½ watt	(brown-green-red)
 CR7	1N914/4820-0201	Check for orientation!
 R26, R27	$27k\Omega$ $\frac{1}{4}$ watt	(red-purple-orange)

	R28	240Ω ½ watt	(red-yellow-brown)
	R31-R34	1kΩ ¼ watt	(brown-black-red)
	CR11	1N4585	Check for orientation!
	Sockets Z1, Z	Z2, Z5	Pin #1 to lower left.
	Socket Z4		Pin #1 to lower left.
	Sockets Z6, Z Z17, Z18	Z7, Z11, Z12, Z16,	Pin #1 to lower left.
***************************************	Sockets Z15, Z22, Z23	Z9, Z20, Z21, 3, Z24	Pin #1 to lower left.
	Sockets Z3, Z	z8, Z19	Pin #1 to lower left.
-	Sockets Z13,	Z14, Z19	Pin #1 to lower left.
	Connector J1		
	Socket/connec	ctor J2	
***********	Socket/connec	tor J3	
-	C9, C10, C11	39 μfd tantulum	Check for orientation!
	R29, R30	470Ω 1 watt	(yellow-purple-brown)
	C27, C28, C7	by-pass capacitors	
	C1-C3	.01 μfd blocking capacitors	
	C4	.1 $\mu$ fd filter capacitor	
	C6	.001 $\mu$ fd blocking capacitor	
	C12-C15	by-pass capacitors	
	C16-C19	by-pass capacitors	
	C20-C26	by-pass capacitors	
earniterio vicino	and hand tigh and board. S	M340.5/7805 5 v. regulator by tening the nut and bolt throug older the leads. If heat sink and the board, heat sink and the board, heat sink and the board,	h the regulator, heat sink, grease is available, apply
	Q1	2N3904 transistor	Check for orientation!
	Q2-Q4	2N3906 transistor	Check for orientation!
	K1-K3	761A05 Form A Gordos relay	Check for orientation!



#### POWER-UP AND SYSTEM CHECK OUT

## POWER SUPPLY/VOLTAGE REGULATOR CHECK OUT

Voltage requiremen	its:	(reference to grou	nd - pins 50 and 100)
Pins 1 and 51		than 7 volts than 10 volts	approx7 amps
Pin 2		than 14 volts than 22 volts	approx1 amps
Pin 52		than -22 volts than -14 volts	approx1 amps

Before installing any of the integrated circuits, apply power to pins 1 and 51, pin 2, and pin 52 (ground at pins 50 and 100) as specified above. Power can come from external supplies or from the host computer. If the board is to be plugged into the computer, it is advisable to use an extender board during testing of the interface. After applying power, perform the following measurements with a volt meter:

- (1) J3 pin 4 -12 volts
- (2) Z5 pin 11 +12 volts
- (3) Z1 pin 16 +5 volts
- (4) Z16 pin 16 +5 volts
- (5) Z15 pin 16 +5 volts
- (6) Z8 pin 14 +5 volts
- (7) Z24 pin 16 +5 volts

If the voltage at any of the check points differs by more than 5% of the required value, return the board for trouble shooting and repair.

## POWER-UP CHECK OUT

Install the integrated circuits as per the layout sheet. Be very careful not to bend pins of the ICs under the package. Often a pin which is bent under an IC appears to be inserted in the socket. REMEMBER: BENT PINS ARE THE MOST COMMON REASON FOR A MALFUNCTIONING BOARD!

After the parts have been installed, repeat the seven voltage measurements specified in the previous section. If the test points differ from the required values, the board should be returned. Under no circumstance should sockets be removed from the board. Such abuse may void the warranty under which repair and service is performed.

## SYSTEM CHECK-OUT INTRODUCTION

The I/O interface occupies 1024 bytes of memory: 200:000 through 203:377 octal or 8000 through 83FF hex. It also uses I/O device addresses 4, 5, 6 and 7. These addressing requirements can be easily customized for special applications but in the check-out procedures which follow, the standard addressing will be used. Users of customized boards should apply appropriate address offsets in the test programs below.

The 1024 bytes of memory used by the I/O board are evenly divided between RAM and PROM. The PROM occupies the first 512 bytes and the RAM uses the last 512 bytes. The board uses I/O device 4 for the cassette data input, data output, and tape motion control. Device 5 is dedicated to the serial data port which includes both a TTY and an RS232 interface. Device 7 controls the 8 bit bi-directional parallel port. Device 6 performs two distinct tasks: the output half controls the paper tape reader, if one is present, on the TTY connected to device 5. The input half furnishes data ready status bits for both the serial port and the parallel port.

The check out procedures which follow are designed to exercise each I/O device on the board along with the RAM and PROM. The RAM is also utilized to facilitate the task of checking the I/O devices.

At a number of places in the check out procedure, it is necessary to set sense switches and start short test programs. For IMSAI or ALTAIR computers a program is started by doing an "examine" at the starting address and then pressing the RUN switch. For the EQUINOX computer or a Keyed-Up 8080 Front Panel/CPU, a program is started by initializing the Program Counter of the CPU and pressing the M key. For computers which have no formal front panel such as the SOL or VECTOR GRAPHICS, a test program is started by having the system monitor perform an "EXECUTE" command to the test program's starting address.

There are several short 28 gauge wires included in I/O Bag #1. These wires are to be used to jumper pins J3 in the check out tests that follow. These wires are fine enough and stiff enough so that they may be inserted into the various connectors of the connector/socket J3.

## ROM

Insert the I/O board in the main frame and turn on the computer. After power up, perform the following examine memory contents:

0cta	1	Hex	
200:000	303	8000	С3
200:001	003	8001	03
200:002	200	8002	80
200:003	075	8003	3D

201:012	107	810A	47
201:013	042	810B	22
201:014	336	810C	DE
201:015	203	810D	83

If there are any deviations from the above values, the board should be returned for service. This completes the ROM check out.

## RAM

Using the computer's front panel or the system monitor, enter the following data starting at address 202:000 octal (8200 hex):

0ct	al	Hex		
202:000	333	8200	DB	START IN SSW
202:001	377	8201	FF	
202:002	323	8202	D3	OUT TAPE
202:003	004	8203	04	
202:004	303	8204	C3	JMP START
202:005	000	8205	00	
202:006	202	8206	82	

After the data has been entered, examine the locations to be sure the RAM has accepted it correctly. This program will be used along with the sense switches (front panel numeric switches on the EQUINOX or Keyed-Up 8080 panel) to check output port 4, the tape cassette command and data output port. If the RAM will not accept data correctly, return the board for service.

Again, using the system monitor or the computer's front panel, enter the following data starting at address 203:000 octal (8300 hex):

BEGIN IN SSW	DB	8300	333	203:000
	FF	8301	377	203:001
OUT PARALLEL	D3	8302	323	203:002
	07	8303	007	203:003
IN PARALLEL	DB	8304	333	203:004
	07	8305	007	203:005

203:006	062	8306	32	STA TEST
203:007	014	8307	ОС	
203:010	203	8308	83	
203:011	303	8309	С3	JMP BEGIN
203:012	000	830A	00	
203:013	203	830B	83	
203:014	000	830C	00	TEST

As before, examine the data to be sure that the RAM has accepted it correctly. If there are memory cells which do not accept data, return the board.

This second sequence of data is another program which will help in checking out the input and output functions of the parallel port -- port 7. To check out the various I/O ports, several items are necessary:

- 1. Some jumpers at least three inches long (included in Bag #1).
- 2. Ohm meter.
- 3. Voltmeter or logic probe.

## TAPE COMMAND AND DATA OUT PORT: I/O DEVICE 4

The first short program entered at 202:000 octal (8200 hex) will be used; it is recapped here for convenience.

0ctal			×	
202:000	333 377	8200	DB FF	START IN SSW
202:002	323 004	8202	D3 04	OUT 4
202:004	303 000 202	8204	C3 00 82	JMP START

Start this program with all the sense switches off. (The sense switches, i.e., the numeric keys of the EQUINOX's front panel are automatically reset whenever the front panel program transfers control to the user's program by pressing the M switch.)

Using the volt meter or logic probe, measure the voltage at the left side of R7 (the 7th resistor below J2). The voltage should be almost zero (TTL logic zero).

Turn on sense switch 0 (press numeric pad switch 0 on the EQUINOX or Keyed-Up 8080). Verify that the voltage is now approximately 4 volts (TTL logic 1). Turn sense switch 0 off again (press S, then M on the EQUINOX).

Verify that the level at the left side of R7 is again a logic zero. Stop the program. This completes the check out of the tape cassette data output.

Restart the program. Connect an ohm meter across pins 6 and 11 of J2. Verify that this is an open circuit. Also, verify that the logic level at pin 12 of Z8 (7406) is a logic 0. Turn on sense switch 1. Verify that pins 6 and 11 of J2 now make a short circuit and that the level at pin 12 of Z8 is now 5 volts. Turn off sense switch 1. Verify that pins 6 and 11 of J2 are again open circuited while the level at pin 12 of Z8 is again zero. This completes the check out of the unit #1 tape command port -- Z8 enables data input from the tape channel 1 and pins 6 and 11 connect to the motor control relay at channel 1.

Connect the ohm meter next across pins 4 and 12 of J2 and verify that this is an open circuit. Also, verify that pin 6 of Z8 is a logic 0. Turn on sense switch 2. Check that pins 4 and 12 of J2 are now short circuited and that pin 6 of Z8 is a logic 1 (5 volts). Turn off sense switch 2 and check that pin 6 of Z8 has returned to a logic zero (approximately 0 volts) and that pins 4 and 12 of J2 are again an open circuit. This completes the check out of the unit #2 command port.

Finally, connect the ohm meter across pins 2 and 15 of J2 and check that this is an open circuit. Also check that pin 10 of Z8 is at a logic 0. Turn on sense switch 3 and check that pins 2 and 15 of J2 are now shorted while pin 10 of Z8 is at 5 volts. Turn sense switch 3 off now and check that at open circuit condition is again maintained across pins 2 and 15 of J2 while the logic level at pin 10 of Z8 has returned to zero. This completes the checkout of the tape output data and command port. Should any of the above tests fail, please return board for service.

## TAPE INPUT DATA PORT:

Enter the following program at 202:000 octal (8200 hex):

202:000	333 004	8200	DB 04	GET IN 4
202:002	346 001	8202	E6 01	ANI 1
202:004	062 012 202	8204	32 OA 82	STA DATA
202:007	303 000 202	8207	C3 00 82	JMP GET
202:012	377	820A	FF	DATA 377

Remove the LM311 from socket Z5A. Connect a jumper from pin #15 of Z5 to pin 3 of J3. Start the above program. Stop the program and examine location 202:012. Verify that this location contains zero. Next, jumper pin 15 of Z5 to pin 4 of J3. Restart the above program. Stop the program and verify that location 202:012 now contains the value 001. Remove the jumper and replace the LM311 in Z5A.

This completes the check out of the tape data input channel. The test does not exercise the LM311 comparator or its input circuitry but it does verify that the output from the LM311 is being properly accessed by the CPU. Later, we will exercise the input circuitry of the tape data channel.

## SERIAL INPUT DATA PORT:

Enter t	he fol	lowing	program a	t 202:000	octal	(8200	hex):
---------	--------	--------	-----------	-----------	-------	-------	-------

202:000	333 005	8200	DB 05	GETS IN 5
202:002	346 001	8202	E6 01	ANI 1
202:004	062 021 202	8204	32 11 82	STA SDATA
202:007	333 006	8207	DB 06	IN 6
202:011	346 001	8209	E6 01	ANI 1
202:013	062 022 202	820B	32 12 82	STA STATUS
202:016	303 000 202	820E	C3 00 82	JMP GETS
202:021	377	8211	FF	SDATA 377
202:022	376	8212	FE	STATUS 376

After entering any data into a computer, especially a test program such as the one above, be sure to examine the contents for accuracy. After the program has been checked, start the test. Stop the program and examine locations 202:021 (8211 hex) and 202:022 (8212 hex). Location 202:021 should now contain 000 and location 202:022 should contain 001.

Using a jumper, connect pin #4 to pin #5 of J3. Restart the test program. Stop the program and examine locations 202:021 and 202:022 for values 001 and 000, respectively. Disconnect the jumper. This completes the serial input data port check out.

## SERIAL OUTPUT DATA PORT:

Enter the following program at 202:000 octal (8200 hex):

00	ΓAL		HEX	
202:000	333 377	8200	DB FF	TEST IN SW
202:002	323 005	8202	D3 05	OUT 5
202:004	303 000 202	8204	C3 00 82	JMP TEST

Using a jumper, connect pins 6 and 7 of J3 together. Make sure all the sense switches are off (the sense switches are always turned off when the front panel of the EQUINOX is active). Start the program. With a volt meter, verify that the voltage between pin 6 of J3 and ground is 0 volts and that the voltage between pin #2 of J3 and ground is +12 volts.

Turn on sense switch 0 (press the 0 digit on the EQUINOX). Measure the voltage between pin #6 of J3 and ground. It should be 5 volts. Measure again the voltage between pin #2 of J3 and ground. It should now read -12 volts instead of +12 volts. Stop the program and disconnect the jumpers. This completes the test of the serial output data port.

## TTY/PAPER TAPE READER CONTROL PORT:

Enter the following program at 202:000 octal (8200 hex):

202:000	333 377	8200	DB FF	TEST IN SSW
202:002	323 006	8202	D3 06	OUT 6
202:004	333 000 202	8204	C3 00 82	JMP TEST

Using a jumper, connect pins 8 and 9 of J3. Make sure all the sense switches are off. Start the program. With a volt meter, measure the voltage between pin #8 of J3 and ground. The voltage should be -12 volts. Turn on sense switch 0 and again measure the voltage between pin #8 of J3 and ground. It should now be +5 volts. Stop the program and remove the jumper. This completes the check out of the TTY Paper Tape Reader Control port.

## PARALLEL STATUS PORT:

Enter the following program at 202:000 octal (8200 hex):

202:000	333 007	8200	DB 07	IN 7
202:002	333 006	8202	DB 06	LOOP IN 6
202:004	346 002	8204	E6 02	ANI 2
202:006	062 014 202	8206	32 OC 82	STA STATUS
202:011	303 002 202	8209	C3 02 82	JMP LOOP
202:014	377	820C	FF	STATUS 377

Start and stop the program. Examine location 202:014 octal and verify that it contains 000. Start the program again. Locate R36 on the circuit board (just to the right of Z4, a 20 pin IC #74LS273). Using a jumper, ground the TOP lead of R36. Stop the program and examine location 202:014 octal (820C hex). This location should now contain 002. This completes the check out of the Parallel Status port.

## PARALLEL DATA PORT:

Earlier, the following program was entered at location 203:000 octal (8300 hex):

203:000	333 377	8300	DB FF	BEGIN IN SSW
203:002	323 007	8302	D3 07	OUT 7
203:004	333 007	8304	DB 07	IN 7
203:006	062 014 203	8306	32 OC 83	STA TEST
203:011	303 000 203	8309	C3 00 83	JMP BEGIN
203:014	000	830C	00	TEST 000

Verify that the program is still in memory. Jumper pin #5 of J1 to ground Check as follows:

Step 1 Turn off all sense switches.	Step 2 Start the program.	Step 3 Turn on sense switch:	Step 4 Stop the program.	Step 5 Examine location 203:014 (830C hex).	Step 6 Contents should be:
н	11	0	н	11	001 octal (01 hex)
11	11	1	11	П	002 octal (02 hex)
n	11	2	H	П	004 octal (04 hex)
11	H	3	11	<b>11</b>	010 octal (08 hex)
11	11	4	11	ti	020 octal (10 hex)
11	<b>11</b>	5	11	H	040 octal (20 hex)
11	ti	6	н	11	100 octal (40 hex)
11	11	7	н	11	200 octal (80 hex)

This completes the check out of the parallel data port and the function modules of the I/O board. If any of the test data differs from what has been described above, return the board for service.

There is one further test which should be done, i.e., exercise the input circuitry of the tape cassette interface. The best way to do this is by performing the examples which illustrate the use of the tape cassette interface (see the tape cassette section of the Principles of Operation).

#### ROUTINE TO GENERATE FLUX CHANGES ON THE CASSETTE TAPE INTERFACE

#### Calling conventions:

- 1. The H-L pair contains the total number of flux changes desired.
- 2. Register C contains a value which determines the delay time between flux changes given by  $32\cdot(c)$  + 65 machine cycles.
- 3. Register A contains a value which determines the delay time from entry of the routine to the first flux changes given by  $32 \cdot (A) + 23$  machine cycles.

There is a delay of 55 machine cycles from the last flux change to the next instruction of the calling routine which follows the "CALL FLUX" instruction.

For a flux change frequency of 1200 hz, or a zero bit, register C should contain 24 decimal. For a flux change frequency of 2400 hz, or a one bit, register C should contain 11 decimal.

All timing assumes one wait state per fetch.

54	200:000	303 003 200	FLUX	JMP	FLUX+3	TIMING PAD
	003	075		DCR	Α	DECREMENT DELAY COUNT
	004	302 000 200		JNZ	FLUX	TEST FOR EXPIRATION
	007	074		INR	Α	COMPLIMENT BIT Ø OF
	010	250		XRA	В	REG B & LEAVE IN ACC
	011	323 004		OUT	TAPE	CHANGE THE FLUX
	013	107		MOV	B,A	UPDATE REG B
	014	053		DCX	H	DECREMENT FLUX CHANGE
	015	175		MOV	A,L	COUNT & TEST FOR
	016	264		ORA	Н	COUNT EQUAL ZERO
	017	171		MOV	A,C	LOAD A WITH DELAY COUNT
	020	302 000 200		JNZ	FLUX	STATUS REG IS SET BY ORA
	023	311		RET		FLUX CHANGES DONE

COPE

#### WRITE TAPE CASSETTE SUBROUTINE

Routine to write a block of data on a cassette recorder using the cassette interface.

#### Calling conventions:

- 1. The entry point to the routine is WTAPE as opposed to the first instruction which has the label TLOOP.
  - 2. The H-L register pair is the address pointer to the data block.
  - 3. The D-E register pair holds the word count of the data block.
- 4. Register C contains a value which determines the delay time from entry of the routine to the first flux change given by  $32\cdot(C)$  + 125 machine cycles.

There is a delay of 158 machine cycles from the last flux change on the audio cassette interface and the instruction of the calling routine following the "CALL WTAPE" instruction.

#### Register Map:

- A: General purpose and delay counts.
- B: Tape command except for bit Ø, register B is preserved. All outputs to the cassette interface are issued through this register.
- C: Delay constants.
- D-E: Data: word count.
- H-L: Flux change count; address pointer.

Including the return address of the calling program, eight levels of the stack are used by the routine. All delays are calculated on the assumption that there is one wait state encountered per memory fetch (both data and instruction memory).

200:024 303 027 200 TL00P

112 303 054 200

027 267

030 341

031 321

JMP

ORA Α

POP

POP D

TLOOP+3

Н

TIMING PAD

TIMING PAD

CURRENT DATA POINTER

CURRENT WORD COUNT

DELAY CONSTANT

COPE

#### READ TAPE CASSETTE WITH RUNNING SUM START BIT DETECTION AND STORE/VERIFY/MOVE OPTIONS

#### Calling conventions:

- 1. The H-L register pair is loaded with the starting address of the area where the data is to be conditionally stored or verified.
- 2. The D-E register pair is loaded with the length of the data block to be read.
  - 3. Register C determines the mode the read routine will operate in.
- a. If bit 6 of register C is set, the routine will verify data on the tape cassette against the contents of the block of memory determined by the two register pairs H-L and D-E.
- b. If bit 7 is set, the routine will simply advance across the tape the number of bytes determined by D-E. This option allows the user to advance exactly one full record on the tape to facilitate computer controlled searches.
- c. The routine assumes that motion control and channel selection will be performed by the calling program.

The routine has logic which rejects wave forms which are not in the region of 1200 to 2400 hz. As such, the tape mechanism may be engaged after the routine has been called but it is not advisable to do so. The routine will wait for signals for approximately one minute and if none are detected during this interval, the routine terminates and writes FF in SERR (STATUS ERROR = 203:366). This location is not initialized by the on-board software. If the verify option is chosen and there is a byte in memory which does not agree with its counterpart on the cassette, DERR (DATA ERROR = 203:365) is filled with some non-zero value. Again this cell is not initialized by the on-board software. Including the return address of the calling program, ten levels of the stack are used by the routine. All timing assumes one wait state per fetch.

200:115				TREAD	PUSH	В	SAVE OPTIONS
116	325					D	SAVE WORD COUNT
117	345				PUSH	Н	SAVE ADDRESS POINTER
120	041	000	360		LXI	H,ØFØØØH	BAD WAVE FORM COUNT
123	345			SYNC	PUSH	H	SAVE
124	257				XRA	Α	INITIALIZE
125	041	370	203		LXI	H,PNTR	8 MOST
130	021	367	203		LXI	D,TOTAL	RECENT
	022		-		STAX	D	TRANSITION TIME
134	167			STORE	MOV	M,A	MEMORY LOCATION
					INR	L	AND THEIR
		134	200		JNZ	STORE	TOTAL
130 133 134 135	021 022	367	203	STORE	LXI STAX MOV INR	D,TOTAL D M,A L	RECENT TRANSITION TIME MEMORY LOCATION AND THEIR

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200:141	107		MOV	B,A	
142	333 004		IN	TÁPE	
144	117		MOV	C,A	
145	014	SUM	INR	c	
146	333 004		IN	TAPE	REG B
150	251		XRA	C	
151	037		RAR	·	IS A MEASURE OF THE
152	004		INR	В	TIME BETWEEN THE
153	312 341 200				LAST TRANSITION AND THE NEXT
156	322 146 200		JZ	LONG	ERROR IF B OVERFLOWS
			JNC	SUM+1	
161	054		INR	L	ADVANCE TRANSITION
162	076 370		MVI	A,MASK	TIME BUFFER
164	265		ORA	L	POINTER
165	157		MOV	L,A	& UPDATE
166	076. 034		MVI	A,HIGH	PROTECT AGAINST
170	270		CMP	В	LOW FREQUENCY
171	332 341 200		JC	LONG	WAVE FORMS
174	032		LDAX	D	GET OLD RUNNING SUM
175	226		SUB	М	SUBTRACT OLDEST TRANSITION TIME
176	200		ADD	В	ADD NEWEST TRANSITION TIME
177	022		STAX	D	SAVE
200	336 164		SBI	CONST	COMPARE FOR
202	322 216 200		JNC	BITS	
205	160		MOV		A START BIT TOTAL
206	006 003			M,B	SAVE NEWEST TRANSITION TIME
210		D4D1	MVI	B,3	OFF SET COUNT
	303 213 200	PAD1	JMP	PAD1+3	TIMING PAD
213	303 145 200		JMP	SUM	MEASURE NEXT TRANSITION TIME
				- 0	
200:216	036 010 .	BITS	MVI	E,8	NUMBER OF BITS TO COLLECT
220	056 013		MV·I	L,11	OFFSET TO DETERMINE Ø/1 WAVE
222	026 017	RENTR	MVI	D,15	BIT CELL TIME ON TAPE
224	046 014		MVI	H,12	TRANSITION COUNTER
226	014	COUNT	INR	C	MEASURE
227	333 004		IN	TAPE	THE
231	251		XRA	С	NEXT
232	037		RAR		TRANSITION
233	055		DCR	L	TIME
234	322 227 200		JNC	COUNT+1	ON THE TAPE
237	056 015		MVI	L,13	OFFSET FOR TRANSITION COUNTER
241	172		MOV	A,D	GET CURRENT CELL TIME VALUE
242	372 252 200		JM	TWO	
245	326 001		SUI	1	SUBTRACT 1 IF WAVE FORM
247	303 257 200		JMP	NEXT	WAS PART OF A ZERO BIT
252	326 002	TWO	SUI	2	OTHERWISE SUBTRACT 2
252 254		IWO			
	303 257 200	NEVE	JMP	NEXT	
257	127	NEXT	MOV	D,A	UPDATE CELL TIME VALUE
260	045		DCR	Н	DECREMENT TRANSITION COUNTER
261	322 226 200		JNC	COUNT	GO BACK IF CELL TIME NON-ZERO
264	174		MOV	A,H	IF H POSITIVE A ZERO BIT
265	027		RAL		OTHERWISE A ONE BIT

200:266	170		MOV	A,B	SHIFT LATEST
267	037		RAR		BIT INTO
270	107		MOV	B,A	DATA REGISTER
271	055		DCR	Ĺ	ADJUST OFFSET
272	055		DCR	ī	
273	177		MOV	A,A	OF TRANSITION TIME COUNT
274	035		DCR		PADDING
275	302 222 200			E	DECREMENT BIT COUNT
			JNZ	RENTR	
300	361		POP	PSW	DISCARD BAD WAVE FORM COUNT
301	341		POP	H	GET DATA POINTER
302	321		POP	D	GET WORD COUNT
303	361		POP	PSW	GET OPTIONS
304	365		PUSH	PSW	SAVE OPTIONS
305	312 326 200		JΖ	VERFY	TEST FOR VERIFY OPTION
310	332 314 200		JC	INCR	TEST FOR NO LOAD OPTION
313	160		MOV	M,B	LOAD MEMORY WITH DATA
314	043	INCR	INX	Н	ADVANCE DATA POINTER
315	033		DCX	 D	DECREMENT WORD COUNT
316	172		MOV	A,D	TEST FOR
317	263		ORA	E E	WORD COUNT
320	302 116 200		JNZ	TREAD+1	
323	361		POP	PSW	EQUAL ZERO
324	257				DISCARD OPTIONS REGISTER
			XRA	Α	A=Ø MEANS SUCCESSFUL READ
325	311		RET		
326	176	VERFY	MOV	A,M	COMPARE PRESENT
327	250	V-1111 1	XRA	B	DATA WITH MEMORY
330	312 314 200		JZ	INCR	CONDITIONALLY SET DERR
333	062 365 203		STA		
336	303 314 200		JMP	DERR	WITH NON-ZERO VALUE
330	303 314 200		JMP	INCR	RETURN
341	341	LONG	POP	Н	GET BAD WAVE FORM
342	043		INX	Н	COUNT & INCREMENT
343	174		MOV	A,H	TEST
344	265		ORA	L .	COUNT FOR
345	302 123 200		JNZ	SYNC	OVERFLOW
350	075		DĊR	A	OVER FLOW OCCURRED
351	062 366 203		STA	SERR	SET STATUS ERROR
354	341		POP		
355	321			H	RESTORE
356	301		POP	D	STACK
			POP	В	& RETURN
357	311		RET		TO CALLING PROGRAM

#### CASSETTE OPERATING EXECUTIVE (COPE)

#### Calling conventions:

1. Register A is the command register.

Bit Ø is the write/read command with a one signifying a write.

Bit 1 activates channel 1.

Bit 2 activates channel 2.

Bit 3 activates channel 3.

Bit 7 = 1 signifies that the activated channels shall be stopped upon completion of the present command.

Bit 7 = 0 signifies that the selected channels shall continue to facilitate chained sequences of tape commands.

2. Register C is the options register for read commands, specifically:

Bit 6 = 1 will prevent loading of the read data and will perform a compare operation with the tape data and the data pointed to by the H-L and D-E register pairs.

Bit 0 = 1 will prevent loading of the read data and simply move the tape forward the number of bytes specified by the D-E register pair. This option is useful for passing over exactly one record so as to be properly positioned to read the next record.

- 3. Register pair H-L is the data pointer which should be set to the starting address of the block of data that is to be read, written, or verified.
- 4. Register pair D-E is the word or byte count register and should be initialized to the length of the data block.

The routine makes extensive use of the stack and as a result, rather than rely on the user to provide the necessary stack space, the calling program's stack pointer is saved. The local RAM of the interface is used as a temporary stack during the duration of the routine. After completing its operations, COPE restores the calling program's stack pointer.

COPE

201:012	107	COPE	MOV	B.A	INITIALIZE COMMAND REG FOR TAP SAVE ADDRESS POINTER	E
013	042 336 203		SHLD	LSAVE	SAVE ADDRESS POINTER	
016	041 000 000		LXI	H.Ø	GET	
021	071		DAD	SĎ.	STACK POINTER	
022	061 336 203		LXI	SP.LSAVE	SAVE STACK POINTER ON	
025	345		PUSH	Н	COPE STACK	
026	323 004		OUT	TAPF	START CASSELLE	
030	365		PUSH	PSW	SAVE COMMAND	
031	041 000 000		LXI	H,Ø	DO A	
034	053	WCOPE	DCX	H.	DELAY	
035	174		MOV	A,H	OF APPROXIMATELY	
036	265		ORA	L	ONE	
037	302 034 201		JNZ	WCOPE	SECOND	
042	361		POP	PSW	GET COMMAND	
043	365		PUSH		AND SAVE AGAIN	
044	037		RAR		TEST FOR A READ	
045	322 104 201.		INC	READ	OPERATION	
050	041 300 135		1 7 1	H 2 PQQQQ	GENERATE	
053	016 013		MVI	C.11	5 SECONDS OF 2400 HZ WAVE REST BITS DELAY CONSTANT	
055	171		MOV	A.C	OF 2400 HZ WAVE	
056	315 000 200		CALL	FLUX	REST BITS	
061			MVI	0.6	DELAY CONSTANT	
	052 336 203		LHLD	LSAVE	GET THE ADDR POINTER	
066	315 034 200		CALL	WTAPE	WRITE RECORD	
	361	EXIT	POP	PSW	GET COMMAND	
071		LATI	ORA	Α	TEST TO TURN	
072	362 101 201		JP	RETRN	OFF CASSETTE	
	257		XRA	A		
070	323 004		OUT	TAPE	STOP CASSETTE	
101	341	RETRN	POP	H	GET OLD STACK POINTER	
101	371	1,21101	SPHL		RESTORE	
103	311		RET		RETURN	
נטו	110		/ <b></b> /			
104	052 336 203	READ	LHLD	LSAVE	GET ADDRESS POINTER	
	315 115 200		CALL			
112	303 071 201		JMP	EXIT		
112	JUJ U/1 ZU1		U. II			

COPE

201.115 265

#### COMPUTE CHECK-SUM ROUTINE

#### Calling conventions:

- 1. The register pair H-L is loaded with the starting address of the data block on which the check-sum is to be computed.
- 2. The register pair D-E is loaded with the word count of the data block.
  - 3. The computed check is returned in the register pair H-L.

Including the return address of the calling program, the routine uses four levels of the stack.

201:115	345	CHECK	PUSH	Н	SAVE ADDRESS POINTER
116	041 000 0	000	LXI	H,Ø	INITIALIZE CHECK SUM
121	104		MOV	В,Н	
122	343	GDATA	XTHL	•	SAVE & EXCHG/ADDR POINTER
123	116		MOV	C,M	GET DATA
124	043		INX	H	INCREMENT ADDRESS POINTER
125	343		XTHL		SAVE & GET PARTIAL CHECK SUM
126	011		DAD	В	ADD NEW DATA
127	033		DCX	D	DECREMENT WORD COUNT
130	172		MOV	A,D	TEST FOR
131	263		ORA	E	WORD COUNT
132	302 122 2	01	JNZ	GDATA	EQUAL ZERO
135	321		POP	D	RESTORE STACK
136	311		RET		
137	061 336 2	03 BOOTS	LXI	SP.LSAVE	READ
142	041 000 2	02	LXI	H.MSA+200:000	
145	021 000 0		LXI	D,256	FIRST
150	076 202		MVI	A.2020	RECORD
152	016 000		MVI	C,Ø	AND
154	303 012 2	01	JMP	COPE	BRANCH THERE
_					

COPE

## THE SERIAL DATA INPUT ROUTINE WITH READER CONTROL AND ECHO OPTIONS

### Calling conventions:

- 1. The on-board RAM location SCON (SERIAL SPEED CONSTANT) must be initialized to the proper value, preferably using the DETCT routine below. This is so that the cell time assumed by the routine and the cell time of the device connected to the interface are compatible.
  - 2. Register B is the options register.
- a. If bit  $\emptyset$  is a one, echoes to the printer are suppressed; otherwise, the data will be echoed as collected.
- b. If bit 7 is a one, the paper tape reader, if one is connected to the interface, will be turned on until a start bit is encountered and then turned off.
  - c. The routine returns with the collected data byte in register D.

including the return address of the calling program, six levels of the stack are used by the routine.

#### Register Map:

- A: General purpose accumulator.
- B: The data collection options register described above.
- C: Not used.
- D-E: 16 bit storage shift register for data collection.
- H-L: Delay time counts.

All timing assumes one wait state per fetch.

163 166 167 171 173	170 007 323 006 052 363 203 345 036 377 333 005 037	SWAIT SLOOK	MOV RLC OUT LHLD PUSH MVI IN RAR	A,B  READR SCON H E,-1 SERAL	CONDITIONALLY TURN ON THE PAPER TAPE READER GET THE SPEED CONSTANT & SAVE ON THE STACK INITIALIZE ½ THE SHIFT REG GET INPUT DATA & TEST
174	332 171 201		JC	SL00K	FOR A START BIT

201:177	315 240 201		CALL	DELAY	WAIT HALF A
202	341		POP	H	BIT TIME
203	333 005		IN	SERAL	& VERIFY
205	037		RAR		THAT A START
206	332 166 201		JC	SWAIT	BIT IS PRESENT
211	026 377		MV I	D,-1	INITIALIZE OTHER ½ OF SHIFT REG
213	257		XRA	Α	STOP THE
214	323 006		OUT	READR	READER
216	345	GTBIT	PUSH	H	UPDATE STACK
217	051		DAD	Н	CALCULATE SPEED
220	053		DCX	Н	CONSTANT FOR A FULL BIT TIME
221	333 005		IN	SERAL	GET INPUT
223	137		MOV	E,A	UPDATE SHIFT REG
224	315 240 201		CALL	DELAY	DELAY 1 BIT TIME & SHIFT D-E
227	341		POP	Н	GET SPEED CONSTANT
230	332 216 201		JC	GTBIT	WAS START BIT SHIFTED TO CARRY?
233	076 001		MVI	A,1	YES. LEAVE WITH
235	323 005		OUT	SERAL	PRINTER IN SPACE MADE
237	311		RET		

# SERIAL DELAY ROUTINE USED IN CONJUNCTION WITH THE SERIAL INPUT AND OUTPUT ROUTINES

#### Calling conventions:

- 1. H-L is initialized with a value which determines the time before the routine returns to the calling program given by 53(H-L) + 58 machine cycles.
- 2. Bits  $\emptyset$  of registers B and E are initialized to values consistent with these bits being ored together and sent to the serial output device.

#### All timing assumes one wait state per fetch.

201:240	173	DELAY	MOV	A,E	CONDITIONALLY
241	260		ORA	В	OUTPUT BIT Ø OF E
242	323 005		OUT	SERAL	TO THE SERIAL DEVICE
244	053		DCX	н	DECREMENT
245	174		MOV	A,H	H-L PAIR
246	265		ORA	L	AND TEST
247	302 240 201		JNZ	DELAY	FOR ZERO
252	172		MOV	A,D	ROTATE D-E
253	037		RAR		ONE
254	173		MOV	A,E	BIT
255	037		RAR		POSITION
256	137		MOV	E,A	TO THE
257	172		MOV	A,D	RIGHT
260	037		RAR	•	WITH END AROUND
261	127		MOV	D.A	BIT PRESERVED
262	311		RET	-	

COPE

#### SERIAL OUTPUT ROUTINE

#### Calling conventions:

- 1. Register A is initialized to an 8-bit value to be serially sent to the serial output device.
- 2. The on-board RAM location SCON (SERIAL SPEED CONSTANT) must be initialized to the proper value preferably using the DETCT routine below. This is so that the cell time assumed by the routine and the cell time for the device connected to the interface are compatible.

The routine starts by sending a zero for one cell time to the serial device connected to the interface. It next sends bit Ø of register A, followed by bits 1 through 7. It then sends a one for two cell times as rest bits and returns to the calling program. All timing assumes one wait state per fetch. Including the return address of the calling program, four levels of the stack are used by the routine.

#### Register Map:

- A: Serial output data.
- B: The routine loads B with twice A to force output when the delay routine is called.
  - C: Used as a bit counter and initialized to 11 decimal.
- $\ensuremath{\text{D-E}}\colon$  16-bit storage shift register for output to the serial output device.

H-L: Delay time counts.

201:263	207	SROUT	ADD	Α	ADD A START BIT
264	107		MOV	В,А	MAKE BIT $\emptyset$ OF B = $\emptyset$
265	137		MOV	E,A	SHIFTED DATA TO E
266	076 013		MVI	A,11	THIS IS BIT COUNT & REST BITS
270	117		MOV	C,A	COUNT TO REG C
271	027		RAL		LOAD D WITH REST BITS
272	127		MOV	D,A	& HIGH ORDER DATA BIT
273	052 363 203	OLOOP	LHLD	SCON	GET SPEED CONSTANT
276	052 363 203		LHLD	SCON	PADDING
301	051		DAD	Н	ADJUST FOR OUTPUT
302	053		DCX	Н	L00P
303	315 240 201		CALL	DELAY	OUTPUT DATA BIT & SHIFT
306	015		DCR	С	DECREMENT BIT COUNT
307	302 273 201		JNZ	OLOOP	
312	311		RET		
• -					

#### SERIAL DEVICE SPEED DETECTION ROUTINE

This routine should be used when the host system is first turned on. The serial input device that will be communicating with the interface should be connected to the interface but should not be turned on until just before the routine is started. The starting address of the routine 201:313 (81CB hex) should be set on the switches followed by an examine. The serial device should then be turned on and the routine started directly afterward. Depending upon certain initial conditions in the host system, the interface itself, and the serial device, the printing mechanism of the serial output device may momentarily not be initialized correctly. However, the first task of the DETCT routine after it initializes the stack pointer is to make sure that the output to the serial printer is correct in accordance with the input from the serial keyboard device.

After these steps are accomplished, the user should then strike the return or carriage return key on the keyboard. The DETCT routine is expecting to see this ASCII character and when it does, it is possible for the routine to measure the baud rate of the device and initialize the onboard RAM location SCON (SERIAL SPEED CONSTANT). After accomplishing this task, the routine comes to a dynamic halt (JMP\*) and the user can then halt the host system and start the COPE BOOTSTRAP routine.

All timing assumes one wait state per fetch.

6	201:313	061 340 20	}	LXI	6,MSA+203:340	INITIALIZE STACK POINTER
	316	026 006	DETCT		D,6	TRANSITION COUNTER
	320	112		MOV	C,D	BIT Ø OF C IS
	321	014	TT00P	INR	Ċ .	IMPORTANT BIT HERE
	322	041 001 000	)	LXI	H,1	OFFSET FOR H-L
	325	333 005	GETIN	IN	SERAL	GET INPUT
	327	323 005		OUT	SERAL	ECHO
	331	251		XRA	C	COMPARE WITH BIT Ø OF C
	332	037		RAR		
	333	043		INX	Н	
	334	322 325 201		JNC	GETIN	WAIT IF NO TRANSITION
	337	345		PUSH	Н	SAVE DELAY TIME
	340	025		DCR	D	· -
	341	302 321 201		JNZ	TTOOP	
	344	341		POP	Н	GET 3 ZERO BITS TIME
	345	341		POP	Н	GET 2 ONE BITS TIME
	346	301		POP	В	GET 1 ZERO BIT TIME
	347	011		DAD	В	ADD TOGETHER
	350	301		POP	В	GET 1 ONE BIT TIME
	351	011		DAD	В	ADD TO PREVIOUS SUBTOTAL
	352	301		POP	В	DISCARD START BIT
	353	301		POP	В	& INFINITE STRING
	354	026 003		MVI	D,3	OF REST BITS

#### SERIAL DEVICE SPEED DETECTION ROUTINE

This routine should be used when the host system is first turned on. The serial input device that will be communicating with the interface should be connected to the interface but should not be turned on until just before the routine is started. The starting address of the routine 201:313 (81CB hex) should be set on the switches followed by an examine. The serial device should then be turned on and the routine started directly afterward. Depending upon certain initial conditions in the host system, the interface itself, and the serial device, the printing mechanism of the serial output device may momentarily not be initialized correctly. However, the first task of the DETCT routine after it initializes the stack pointer is to make sure that the output to the serial printer is correct in accordance with the input from the serial keyboard device.

After these steps are accomplished, the user should then strike the return or carriage return key on the keyboard. The DETCT routine is expecting to see this ASCII character and when it does, it is possible for the routine to measure the baud rate of the device and initialize the onboard RAM location SCON (SERIAL SPEED CONSTANT). After accomplishing this task, the routine comes to a dynamic (JMP\*) halt and the user can then halt the host system and start the COPE BOOTSTRAP routine.

All timing assumes one wait state per fetch.

201.212 001 250 000

201:			340			LXI	6, MSA+203:34	OINITIALIZE STACK POINTER
	316	026	006		DETCT	MVI	D,6	TRANSITION COUNTER
	320	112				MOV	C,D	BIT Ø OF C IS
	321	014			TT00P	INR	C	IMPORTANT BIT HERE
	322	041	001	000		LXI	H <sub>2</sub> 1	OFFSET FOR H-L
	325	333	005		GETIN	IN	SERAL	GET INPUT
	327	323	005			OUT	SERAL	ECHO
	331	251	_			XRA.	C	COMPARE WITH BIT Ø OF C
	332	037				RAR	•	COMPARE WITH BIT MOF C
	333	043				INX	н	
	334	_	325	201		JNC	GETIN	WALT IS NO TRANSPICE
	337	345	,			PUSH	H	WAIT IF NO TRANSITION
	340	025				DCR	D	SAVE DELAY TIME
	341		321	201		JNZ	TTOOP	
	344	341	,,,,	201		POP	H	AFE A 4554 AVE
	345	341				POP		GET 3 ZERO BITS TIME
	346	301					H	GET 2 ONE BITS TIME
	347	011				POP	В	GET 1 ZERO BIT TIME
	350					DAD	В	ADD TOGETHER
		301				POP	В	GET 1 ONE BIT TIME
	351	011				DAD	В	ADD TO PREVIOUS SUBTOTAL
	352	301				POP	В	DISCARD START BIT
	353	301				POP	В	& INFINITE STRING
3	354	026	003			MVI	D,3	OF REST BITS

COPE

Hex Listing

## ROUTINE TO GENERATE FLUX CHANGES ON THE CASSETTE TAPE INTERFACE

(see octal listing for comments)

8000	С3	03	80	FLUX	JMP	FLUX+3
8003	3Ď	-			DCR	Α
8004	C2	00	80		JNZ	FLUX
8007	3 C				INR	Α
8008	A8				XRA	В
8009	D3	04			OUT	TAPE
800B	47				MOV	B,A
800C	2B				DCX	Н
800D	7D				MOV	A,L
800E	В4				ORA	Н
800F	79				MOV	A,C
8010	C2	00	80		JNZ	FLUX
8013	С9				RET	

#### WRITE TAPE CASSETTE SUBROUTINE

8014	С3	17	80	TL00P	JMP	TLOOP+3
8017	В7				ORA	A
8018	Ε1				POP	Н
8019	D1				POP	D
801A	0E	12			MVI	C,18
801C	7A			WTAPE	MOV	A,D
801D	В3				ORA	E
801E	C8				RZ ·	
801F	1B				DCX	D
8020	D5				PUSH	D
8021	5E				MOV	E,M
8022	23				INX	Н
8023	E5				PUSH	Н
8024	79				MOV	A,C
8025	16	03			MVI	D,3
8027	21	08	00	ZBIT	LXI	н,8
802A	0E	18			MVI	C,24
802C	CD	00	80	WRITE	CALL	FLUX
802 F	7B				MOV	A,E
8030	B2				ORA	D
8031	CA	14	80		JZ	TL00P
8034	7A				MOV	A,D
8035	1F				RAR	
8036	57				MOV	D,A
8037	7B				MOV	A,E
8038	1F				RAR	
8039	5F				MOV	E,A

803A	3E 00	MVI	A,Ø
803C	7F	MOV	A,A
803D	3E 14	MV I	A,20
803F	D2 27 80	JNC	ZBIT
8042	21 10 00	LXI	Н,16
8045	OE OB	MVI	C,11
8047	7F	MOV	A,A
8048	3E 06	MVI	A,6
804A	C3 2C 80	JMP	WRITE

# READ TAPE CASSETTE WITH RUNNING SUM START BIT DETECTION AND STORE/VERIFY/MOVE OPTIONS

804D	C5			TREAD	PUSH	В
804E	D5				PUSH	D
804F	E5				PUSH	Н
8050	21	00	F0		LXI	H,0F000H
8053	E5			SYNC	PUSH	H i
8054	ΑF				XRA	Α
8055	21	F8	83		LXI	H,PNTR
8058	11	F7	83		LXI	D,TOTAL
805B	12				STAX	D
805C	77			STORE	MOV	M,A
805D	2 C				INR	L
805E	C2	5C	80		JNZ	STORE
8061	47				MOV	B,A
8062	DB	04			IN	TAPE
8064	4F				MOV	C,A
8065	OC			SUM	INR	C
8066	DB	04			IN	TAPE
8068	Α9				XRA	С
8069	1F				RAR	
806A	04				INR	В
806B	ÇA	E1	80		JZ	LONG
806E	D2	66	80		JNC	SUM+1
8071	2 C				INR	L
8072	3E	F8			MVI	A,MASK
8074	В5				ORA	L
8075	6F				MOV	L,A
8076	3E	10			MVI	A,HIGH
8078	в8				CMP	В
807 <del>9</del>	DA	E 1	80		JC	LONG
807C	1A				LDAX	D
807D	96				SUB	M
807E	80		,		ADD	В
807F	12				STAX	D
8080	DE	74			SBI	CONST
8082	D2	ŘΕ.	80		JNC	BITS

8085	70				MOV	M,B
8086	06	03			MVI	В,3
8088	С3	8B	80	PAD1	JMP	PAD1+3
808B	C3	65	80		JMP	SUM
808E	1E	08	••	BITS	MVI	E,8
8090	2E			D113	MVI	L,11
8092	16	0F		RENTR		
8094	26	00		KENIK	MVI	D,15
		UC			MVI	H,12
8096	00	- 1		COUNT	INR	C
8097	DB	04			IN	TAPE
8099	A9				XRA	C
809A	1F				RAR	
809B	2 D		_		DCR	L
809C	D2	97	80		JNC	COUNT+1
809F	2E	OD			MVI	L,13
80A1	7A				MOV	A,D
80A2	FA	AΑ	80		JM	TWO
80A5	D6	01			SUI	1
80A7	С3	ΑF	80		JMP	NEXT
80AA	D6	02		TWO	SUI	2
80AC	С3	ΑF	80		JMP	NEXT
80AF	57	• • • •	••	NEXT	MOV	D,A
80B0	25			116/11	DCR	H .
80B1	D2	96	80		JNC	COUNT
80B4	7C	,,	••		MOV	A,H
80B5	17				RAL	۸,11
80B6	78				MOV	A,B
80B7	1F				RAR	Α,υ
80B8	47				MOV	ВΛ
80B9	2D					B,A
80BA	2D				DCR	L
					DCR	L
80BB	7F				MOV	A,A
80BC	1D		00		DCR	E
80BD	C2	92	80		JNZ	RENTR
80C0	F1				POP	PSW
80C1	E1				POP	Н
80C2	D1				POP	D
80C3	F1				POP	PSW
80C4	F5				PUSH	PSW
80C5	CA	D6	80		JZ	VERFY
80C8	DΑ	CC	80		JC	INCR
80CB	70				MOV	M.B
80CC	23			INCR	INX	H
80CD	1B				DCX	D
80CE	7A				MOV	A,D
80CF	В3				ORA	E
80D0	C2	4E	80		JNZ	TREAD+1
80D3	F1				POP	PSW
80D4	AF				XRA	A
80D5	Ç9				RET	
2007	رب				NE I	

80D6 80D7 80D8 80DB 80DE	7E A8 CA CC 8 32 F5 8 C3 CC 8	3	MOV XRA JZ STA JMP	A,M B INC DER INC
80E1	E1	LONG	POP	H ·
80E2	23		INX	Н
80E3	7C		MOV	A,H
80E4	B5		ORA	L
80E5	C2 53 8	0	JNZ	SYN
80E8	3D		DCR	Α
80E9	32 F6 8	3	STA	SER
80EC	E1		POP	. Н
80ED	D1		POP	D
80EE	C1		POP	В
80E F	C9		RET	

#### CASSETTE OPERATING EXECUTIVE (COPE)

810A	47			COPE	MOV	B,A
810B	22	DE	83		SHLD	LSAVE
820E	21	00	00		LXI	н,о
8111	39				DAD	SP
8112	31	DE	83		LXI	SP, LSAVE
8115	E5				PUSH	Н
8116	D3	04			OUT	TAPE
8118	F5				PUSH	PSW
8119	21	00	00		LXI	H,0
811C	2B			WCOPE	DCX	H
811D	7C				MOV	A',H
811E	B5				ORA	L
811F	C2	10	81		JNZ	WCOPE
8122	F1				POP	PSW
8123	F5				PUSH	PSW
8124	1F				RAR	
8125	D2	44	81		JNC	READ
8128	21	CO	5D		LXI	н,24000
812B	0E	0В			MVI	C,11
812D	79				MOV	A,C
812E	CD	00	80		CALL	FLUX
8131	0E	06			MVI	С,6
8133	2A	DÉ	83		LHLD	LSAVE
8136	CD	10	80		CALL	WTAPE
8139	F1			EXIT	POP	PSW
813A	В7				ORA	A
813B	F2	41	81		JP	RETRN
813E	ΑF				XRA	Α
813F	D3	04			OUT	TAPE
8141	E1			RETRN	POP	Н
8142	F9				SPHL	
8143	С9				RET	

8144	2A	DE	83	READ	LHLD	LSAVE
8147	CD	4D	80		CALL	TREAD
814A	C3	39	81		JMP	EXIT

#### COMPUTE CHECK-SUM ROUTINE

814D	E5			CHECK	PUSH	Н
814E	21	00	00		LXI	н,о
8151	44				MOV	в,н
8152	E3			GDATA	XTHL	
8153	4E				MOV	C,M
	23				INX	H
8155					XTHL	
8156					DAD	В
8157	1B				DCX	D
8158					MOV:	A,D
8159	B3				ORA	E
815A		52	<b>Q</b> 1		JNZ	GDATA
815D		72	٠,		POP	D
					RET	U
815E	C9				NE I	
0.55		~~	02	BOOTS	LXI	CD I CAVE
815F	31		83	BUU 1 3		SP,LSAVE
8162	21				LXI	H,8200H
8165	11	00	01		LXI	D,256
8168		82			MVI	A,82
816A			_		MVI	C,O
816C	C3	0Α	81		JMP	COPE

## THE SERIAL DATA INPUT ROUTINE WITH READER CONTROL AND ECHO OPTIONS

816F	78		INPUT	MOV	A,B
8170	07			RLC	
8171	D3 06			OUT	READI
8173	2A F3	83		LHLD	SCON
8176	E5		SWAIT	PUSH	Н
8177	1E FF			MVI	E,-1
8179	DB 05		SLOOK	IN :	SERA
817B	1F			RAR	
817C	DA 79	81		JC	SLOO
817F	CD AO	81		CALL	DELA'
8182	E1			POP	Н
8183	DB 05			IN	SERA
8185	1F			RAR	
8186	DA 76	81		JC	SWAT
8189	16 FF			MVI	D,-1
818B	ΑF			XRA	Α
818C	D3 06			OUT RE	ADR

818E	E5	GTBIT	PUSH	Н
818F	29		DAD	Н
8190	2B		DCX	H
8191	DB 05		IN	SERAL
8193	5F		MOV	E,A
8194	CD AO 8	1	CALL	DELAY
8197	E1 .		POP	Н
8198	DA 8E 8	1	JC	GTBIT
819B	3E 01		MVI	A,1
819D	D3 05		OUT	SERAL
819F	C9		RET	

# SERIAL DELAY ROUTINE USED IN CONJUNCTION WITH THE SERIAL INPUT AND OUTPUT ROUTINES

81A0	7B		DELAY	MOV	A,E
81A1	BO			ORA	В
81A2	D3 05			OUT	SERAL
81A4	2B			DCX	Н
81A5	7C			MOV	A,H
81A6	B5			ORA	L.
81A7	C2 A0	.81		JNZ	DELAY
81AA	7A			MOV	A.D
81AB	1F			RAR	•
81AC	7B			MOV	A,E
81AD	1F			RAR	•
81AE	5F			MOV	E,A
81AF	7A			MOV	A,D
81B0	1F			RAR	•
81B1	57			MOV	D,A
81B2	C9			RET	•

#### SERIAL OUTPUT ROUTINE

81B3	87	SROUT	ADD	Α
81B4	47		MOV	B,A
81B5	5F		MOV	E.A
81B6	3E 0B		MVI	A,11
81B8	4F		MOV	C.A
81B9	17		RAL	•
81BA	57		MOV	D,A
81BB	2A F3 83	OLOOP.	LHLD	SCON
81BE	2A F3 83		LHLD	SCON
8101	29		DAD	· H
81C2	2B		DCX	Н
8103	CD AO 81		CALL	DELAY
8106	OD		DCR	C
81C7	C2 BB 81		JNZ	OLOOP
81CA	C9		RET	

### SERIAL DEVICE SPEED DETECTION ROUTINE

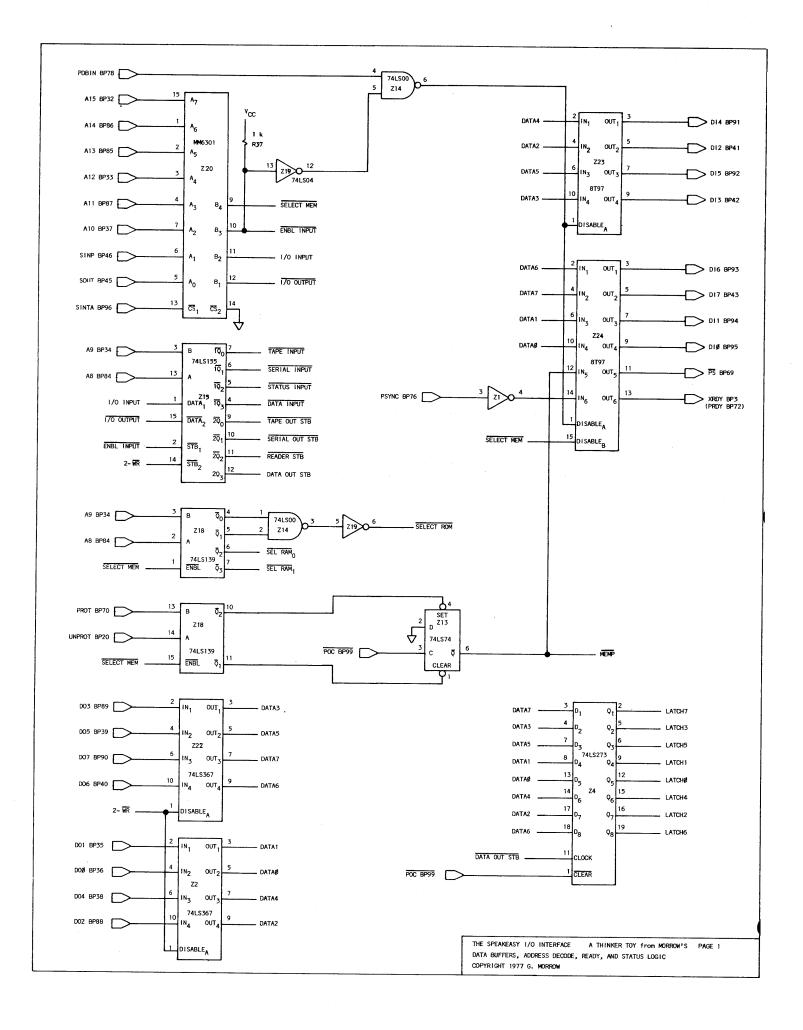
81CB	31 E0	83		LXI	6,83E0
81CE	16 06	-	DETCT	MVI	D,6
81D0	4A			MOV	C,D
81D1	0C		TL00P	INR	C
81D2	21 01	00		LXI	H,1
81D5	DB 05		GETIN	IN	SERAL
81D7	D3 05			OUT	SERAL
81D9	A9			XRA	C
81DA	1F			RAR	
81DB	23			INX	Н
81DC	D2 D5	81		JNC	GETIN
81DF	E5			PUSH	Н
81E0	15			DCR	D
81E1	C2 D1	81		JNZ	TLOOP
81E4	E1			POP	Н
81E5	E1			POP	H
81E6	C1			POP	В
81E7	09			DAD	В
81E8	C1			POP	В
81E9	09			DAD	В
81EA	C1			POP	В
81EB	C1			POP	В
81EC	16 03			MVI	D,3
81EE	AF		SLOOP	XRA	A
81EF	7C			MOV	A,H
81F0	1F			RAR	
81F1	67			MOV	·H,A
81F2	7D			MOV	A,L
81F3	1F			RAR	,-
81F4	6F			MOV	L,A
81F5	15			DCR	D
81F6	C2 EE	81		JNZ	SL00P
81F9	2B			DCX	H
81FA	22 F3	83		SHLD	SCON
81FD		81	STOP	JMP	STOP

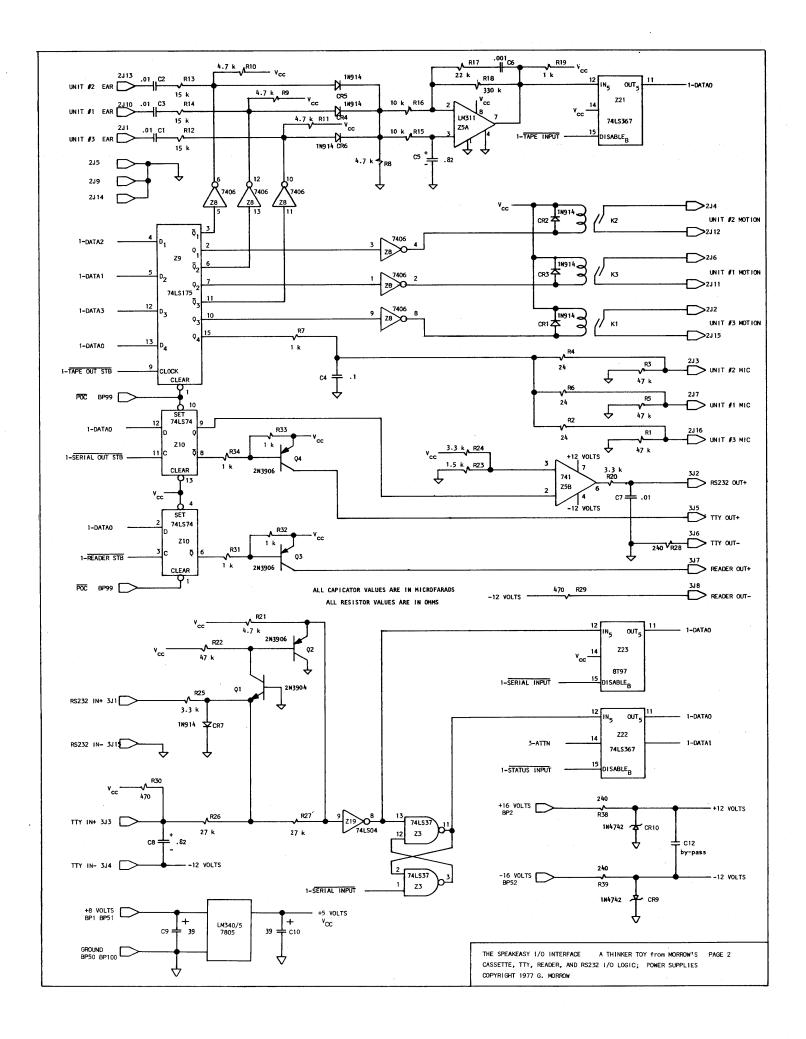
õ

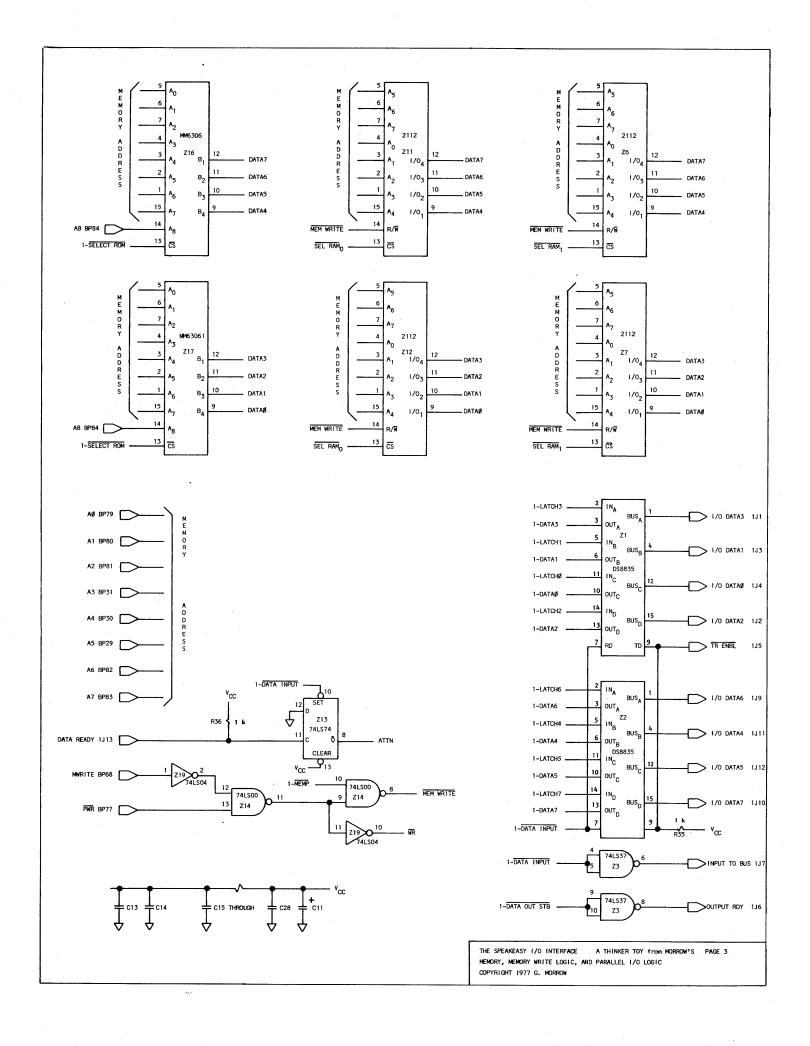
## WARRANTY

Parts are warranted to be free from defects in material and work-manship. Defective parts returned postpaid will be exchanged free of charge. Thinker Toy products purchased in kit form are warranted for six months from date of invoice. Thinker Toy products purchased as assembled units are warranted for one year from invoice date. Malfunctioning units whether purchased in kit form or pre-assembled will be repaired, tested, and returned with a minimal charge for postage/handling if in the opinion of Morrow's Micro-Stuff or Thinker Toys care has been exercised in their assembly and/or use. Warranty is void if on inspection by Morrow's or Thinker Toys it is found that the product has been subject to improper assembly or abuse. Charges will be assessed accordingly for repair parts and labor. Repair fees will not exceed \$25.00 unless prior approval has been obtained from purchaser.

The foregoing warranty is in lieu of all other warranties expressed or implied and in any event is limited to product repair or replacement.







#5347

October 8, 1979

The Speakeasy Inv. 5898

Thinker Toys 5221 Central Ave., #9 Richmond, CA

In accordance with your telephoned instructions, I enclose (1) Rev. 3 of my Z-80 Version of Cope, and (2) a Tape Cassette with recordings made with this version.

Side 1 of the Cassette is a recording from 2K of ROM. When I read this back into RAM, the result was accurate, except that F8 was inserted at the beginning, everything displaced by one byte, and the last byte was lost.

Side 2 of the Cassette is a recording from the above described RAM. When I read this back into another section of RAM, another F8 was inserted at the beginning (making two F8s), everything displaced once more, and another byte lost at the end.

As explained in my Sept. 25 letter, the enclosed program takes care of two things. (1) It addresses Speakeasy on the upper eight address lines in the only way that my computer can (as now configured), and (2) it uses constants that adjust for the lack of WAIT states on each fetch.

I have checked BP3 (XRDY) on the oscilloscope in comparison with the CPU Clock (\$\phi\$). There are clear and distinct signals on Buss 3, but it is a chopped signal and clearly timed wrong for the purpose of creating a WAIT. I am certain that the reason for this is that the circuit on my CPU Board which makes the artificial PSYNC does not (and Can Not) make the signal fast enough to satisfy Speakeasy.

Because of the lack of a WAIT state on each fetch, I am not satisfied that there are not timing problems with the exclosed Z-80 version of COPE. For example, I doubt that the Serial Device Speed Detection will work. I have only tested this program as respects weiting and reading tape.

I hope that you can (1) review the Tape handling portions of the program as respect to needs for more PADS, and (2) the balance in all respects.

The only changes that I have made in this revision are recapped on the next page.

Re Cap of changes made to Cope to produce this Z-80 Version

- 1. INPUT and OUTPUT Macros are used that cause the CPU to address Speakeasy on the upper eight address lines.
- 2. The location of two sub-routines, (1) Compute Checksum, and (2) Boots, are moved ahead of the Executive Routine in order to make room for the IN and OUT macros in the Serial sub-routines.
- 3. Certain Constants are changed so as to compensate for the lack of WAIT states on each fetch.

ORIGINAL	COPE :			
Address		Hex Value	Address:	Comment
801B 802B	12	13 1B	801C 802C	Delay Constant # Flux changes for a Zero bit
803 <b>E</b>	14	16	803E	Set up for a Zero bit
8046	: 0B	OC	8045	Set up flux change for a one bit
80 <b>4</b> 9 8 <b>12</b> 0	: 06 : 0B	07 <b>ф</b> С	8048 812F	Delay Constant 2400 Hz wave rest bits
8132	: 06	07	8135 :	Delay Constant

As stated before, my primary concern is that there may be need for other timing corrections by insering pads. Although there is little room left to do this. There are five NoOps at 80E7 to 80EB, one NoOp at 8154, and two more at the end. The number of machine cycles per line has been entered at the right margin of each page. These assume NO WAIT states.

Sincerely,

Robert 7. Hassard

FA	1	į	2	-80	VERS	JON C	OPE	REV 3	9/22/79
		STRUCTI	ON			SOURCE COD	E		
ADDRESS	BYTE 1	BYTE 2	BYTE 3			OP CODE	OPERAND		#5347
				R	emarks	are	. Same	as in	
				S	PEAKER	8 Y 4	ISER'S	MANUAL	
8000	18	00			FLUX	JR	FLUX+2		12
2	<u>3d</u>					DEC	A FLUX	>28	4
35	20	FB				JR NZ	FLUX	2	12/7
5	30					INC	<u>H</u>		<u>+</u> +
6	A8	<u> </u>			·	XOR	В	: .)	<u> </u>
7	9	4.4				EXX	0404		4
8	01		04	<b> </b>		OUT (c)	0404		12
B	Ed	79	-			COI (c)		l /	4
<u>a</u>	49		ļ			7 4 2	^		<u>T</u>
E	47		<b> </b>	<u> </u>		DEC.	<u> </u>	72	4 6 4
F	2B	ļ	<b> </b>			DEC 1 A D	75	H = 3	<u>0</u>
10	DIL					ORD	<u> </u>	23 5 + 67	+
1	7d B4 79		<del> </del>			EXX LdB DEC LdA ORA LdA	(	1=12 OCH	4
3	20	EB		-		JRN2	FILIX	0=27 1BH	1217
5	<u><u><u></u> <u> </u></u></u>					RET	1 40 %		10
3	<u></u>	1				11 - 1			
		1			WRITE	TAP	E CASS	ETTE	
			<del> </del>		VVI	1 1 1 1 1			
8016	18	00			TLOOP	JR	TLOOP + 2		12
8	27						A		4
9	37 E1 d1					POP	HL		10
A	di					POP	DE	·	10
В	OE	13				Ld C	19 Dec.		7
d	7A				WTAPE	Ld A	D		4
E	<b>B3</b>						E		4
F	C 8					RETZ			11/5
20	18					DEC	DE		6
1	d5					PUSH	DE		
2	5E 23					LdE	(HL)		
3	23	<u> </u>	ļ	<u> </u>	1	INC			6
4	E5	<del> </del>		<b> </b>	<b> </b>	PUSH			
5	79	<u> </u>		<u> </u>	<b></b>	Ld A	C		<u>4</u>
6	16	03	<b></b>	<b> </b>	#	rg D	03		
<b>B</b>	21	08	00	<b> </b>	ZBIT	FGHT	0008		10
	OE		0.5	<del> </del>	1.00175	Ld C	27 Dec.		7
d	Cq	00	80	<del> </del>	WRITE	CHLL	FLUX		1/
30	7B 82	<del> </del>	<del> </del>	-		Ld A,	<del> </del>		<del></del>
	100	EA	<b> </b>	<del> </del>	-	OR JRZ	TLOOP	·	12/7
2	28 7A	E2	+	-	#	110	D		4
5		<del> </del>	<del> </del>	<del> </del>		RRA Ld D,			4
6	57	<del>                                     </del>	1		1	ILA D.	A		4
7	7B	1	<del>                                     </del>			LAA	E		4
8	IF	1	<del>                                     </del>	<b></b>		Ld A, RRA			4
9	5F		<b>†</b>	1	1	LdE	A		4
	·	1	1	1					
<u> </u>		1			<u></u>			:	

,	PA	+GE	2	2-	-80 VER	SION	COPE	REV	3	9/22/79
		STRUCTI			LABEL	OP CODE	OPERAND			
ADDRESS	BA1E 1	BYIE 2	BTIES	BYTE 4			(CONT.)		MMENTS	
803A	3E	00			WKIIE	TAPE LL A	00			7
C	7 <i>F</i>					Ld A	A			4
ď	3E	16				LdA	22 Dec.		····	7
F	30	E7					2 BIT			12/7
41	21	10	00			Ld HL				10
4		OC				rq C		,		71
6		A 7				Ld A	Ą		· · · · · · · · · · · · · · · · · · ·	<u>+</u>
7 9	3E	07 E2				Ld A JR	7 WRITE		**************************************	12
	10	EX				21	WILL	•		
			R	EAD	TAPE	CHSSE	TTE wit	h runnin	g Su	m, etc.
804B	<b>C</b> 5				TREAD	PUSH	BC			11
C	d5					PUSH				11
d	E5					PUSH				11
E	21	00	FO				F000			10
51	E5				SYNC	PUSH				
2	AF	Co	0.3			XOR	A			4
3	21	F8	83				PNTR			10
9	12	F.	83			rg DE				7
Ä	77				STORE					7
В	20			1		INC	L			4
С	20	FC				JR NZ	STORE			12/7
E	47 3E					Ld B	A			4
尸	3E	04				Ld A	04	To Upper	8 Ad	ldr Lines 7
61	dB	04				INA	(04)			
3					SUM	Ld C.	C			
<u> </u>	0C	04			3014	Ld A	04	T- 110000	. 2 AJ	dr Lines 7
7	12	04				INA	(04)	10 OFFE	DH O	11
9	A9					XOR	<u>C</u>			4
A	3 E d B 1 F					RRA				4
<u>B</u>	04					INC	B			4
C	CA	48	80	<u> </u>			LONG	5c		.10
F	30	F4	<b> </b>	ļ		JR NC	SUM + 1			12/7
11	20	0		<u></u>		INC	MASK			4
2 4		F8		<del>                                     </del>		LdA	1.14.216			4
	02	<del>                                     </del>				Ld L	A			4
6	36	1C	<b> </b>	<b></b>		Ld A	HIGH			ול
<b>40%</b>	6F 3E B8					ICP	B			4
9	dA	9 8	80			JMPC	LONG	<u> </u>		10 7
	18	ļ	<u> </u>			Ld 17.				
<u>d</u>	96	<u> </u>				SUB	(HL)	· · · · · · · · · · · · · · · · · · ·	-	7
E	80	<b> </b>				ADDA	1 15			4
7	12					LY (DE	f, 17			
	L	L	1	<del></del>	<u> </u>		1	1		

		Pag	e 3		2-80	UER	SION CO	PE	Rev 3	9/22/79
		STRUCTI	ON			SOURCE COD		_		
ADDRESS	BYTE 1	BYTE 2	BYTE 3	BYTE 4	·	OP CODE	OPERAND	<u> </u>	COMMENTS	
					REAL		PE (CO	<u>ゆナ、)</u>		
8080		74					CONST			4
2	30	07				JR NC	BITS	ļ		12/7
4						79 (Hr)	<u>B</u>			
5	06	03			000	JR 74 B	03	<u> </u>		
7	18	00			PADI	リス	PAD 1 +2	<b>-</b>		12
9	18	d9			BITS	JR	SUM	-	,	12.
B	IE	08			10112	LdE	08	<b></b>		7
d F	2E 16	OF			RENTR	79 D	11 dec			<del></del>
91	26	OC			VENTI	rg H	13 dec			7
3	OC	00			COUNT		- C			4
4		04			COUNT	LdA	04			7
6		04			·	IN A	(04)			11
8	A9					XOR	<u>C</u>	į		4
9	IF					RRA				4
R	2d					DEC	L			4
B	30	F7				JR NC				12/7
d	2E	00				Ld L	13 dec			
F	7A					Ld A,	D			4
HO	FA	A7	80			JMPM				10
3		01				SUB	01	_		4
5	18	04				JR	NEXT			12
7	16	02	<b></b>		TWO	SUB	02			12
9	18	00			NEXT	JR Ld D,	NEXT	-		<del></del>
B	57				NEXI	La V,	H	-		4
<u>C</u>	25	E4				DEC	COUNT			12/7
d F	30 7C	<u> </u>				Ld A	H			4
Bo	17					RLA				4
1	78					Ld A.	В			4
2						LL A.				4
3	47					Ld 13	A			4
3 4 5	20					LAB DEC DEC	L			4
	2d 2d					DEC	<u>L</u>			<u> </u>
6	7F		<u> </u>			Ld A, DEC	A			4
7	1d			<u> </u>		DEC	E	·		4
8 B	C2	8F	80	<u> </u>	1	JMPNZ	RENTR	<del> </del>		10
<u>B</u>	FL	<del>                                     </del>		<b> </b>	<b></b>	POP	AF	1		10
	FI	-	ļ	<del> </del>	<b>H</b> .	POP	HL	-		10
1 0	di	-		<del> </del>	-	POP	DE AF	-	<u> </u>	10
d E F	FL	-	<del> </del>	<del> </del>	-	POP	AC	<del>                                     </del>		
F	F5	od	<del> </del>	-	<b>H</b>	JR Z		1		12/7
CO	28 38	01	<del>                                     </del>	1		JR C	1 NC IS	1		12/2
4	70			<b> </b>	1	LA CHI	INCR DIB			7
5		1	1	1	INCR	INC	Him			4
6		<b>†</b>	1		· · · · · · · · · · · · · · · · · ·	DEC	DE			4
				1						
				7	1			ţ		

		Page	2 4	f 2	-80 VER	SION	COPE	Rev 3	9/22/79
	IN	STRUCTIO	МС			SOURCE COD			
ADDRESS	BYTE 1	BYTE 2	BYTE 3	BYTE 4	LABEL	OP CODE	OPERAND	COMMENTS	
				,	READ	TAI	E (COV	<i>†</i> .)	
8007	7A					Ld A,	D		4
8	<b>B</b> 3					OR 1	E		4
9	<b>C2</b>	40	80			JMPNZ	TREAD+1		10
ς.	FI					POP	AF		10
d	AF					XOR	A		4
E	<b>C9</b>					RET			10
F	7E				VERFY	Ld A	(HL)		7
<u>do</u>					`	XOR	_B		4
	28	FZ				JRZ	ÍNCIS	2550	12/7
3	32	F5	83	:		79 (834	TS), A	DERR	13
6	18	Ed				IK	INCR		12
8	EI				LONG	POP	HL	<u> </u>	10
9	23					INC	HL		4
A	7C					hd 17,	Ħ		4
B	B5		~ -			OR	6		10
<u> </u>	C2	51	80			JMP NZ			
F	34	<b>—</b>	0.5			DEC		6500	13
Eo		F6	83	ļ <u>i</u>		rd (831	76), H	SERR	
3				<b> </b>		POP	HL DE BC		10
- 4				<u> </u>		POP	DE RC		10
5						POP	BC		
6	<u>C9</u>				·	RET			10
	00	00	00						
R	00	00							
	<b></b>			COM	PUTE C	MECK	SUM	ROUTINE	
DAEC		ļ		COM	CHECK	PUSH	14	10011111	11
80EC	E5	00	00		CHECK		0000		10
FO		00	00			Ld B			4
	E3				GDATA	EX (SF	) Hi		19
2					GPILITA	Ld C,			7
2	72					INC	HL		6
4	23 E3	<b> </b>				EX (SP	1.146		19
3 + 5	09					ADD H	BC		15
6						DEC	DE		6
7	7A					Ld A	D		4 4
8						ORA	E		
9	20	F6				JR NZ	G-DATA		12/7
B	11					POP	DE		10
G						RET			10
80Fd	31	dE	83		BOOTS	Ld SP	834E	LSAUE	10
8100	21	00	82			ITG HI	8200		10
3	11	00	01	<b> </b>		Ld DE			10
6	3E	82	<u> </u>			Ld A	82 H		
8	OE	00	<b> </b>	1		ra c	00		
	<b> </b>	<u> </u>	<u> </u>	ļ		JMP	COPE	Continue	. 810A
	<b></b>	<u> </u>	<del> </del>	1					
	1	<u> </u>	<u> </u>	!	,11		1		

PAGE 5					Z-80 V	ERS10	N	COPE		EV	3	9/22/79
INSTRUCTION												
ADDRESS	BYTE 1	BYTE 2	BYTE 3	BYTE 4	LABEL	OP CODE	0	PERAND		COMM	ENTS	
				CA	SSETTE	OPE	RI	TING	EXE	=cu	TIVE	
810A	47	1.0	00		COPE	Ld B					·	4
В	22	dE	83			Tq (831						16
E	21	00	00			ra Hr		000	-			10
	39	1 =	00			ADD HL	. 5	<u> </u>		· · · · · · · · · · · · · · · · · · ·		10
2 5	31	dE	83			LdSP			<u></u>			10
6	E5 d9					PUSH EXX	HL			····		4
7	01	04	n4			79 BC	0	104				10
Á	Ed	79	O I			out (c)	A					12
C	19					EXX						4
त	F5					PUSH	A	<u>C</u>				11
d	21	00	00			LdHL		000				10
21	2B				W COPE	DEC	Н	<b>L</b>				6
2	70					LdA	H					4
2 3 4	<b>B</b> 5					OR	<u></u>					4
4	20	FB				JR NZ		COPE				1217
6	FI					POP	A					10
7	F5					PUSH	A	<u> </u>				11
8	1F					RRA						4
9	30	21				JR NC		EAD				12/7
BE		CO	<u>5d</u>			TG HT	2	1000				10
	OE	OC		<u> </u>		79C		0				
30	19	00 07 dE	00	<b> </b>		Ld A	C					17
1	Ça	00	80	ļ		CALL		LUX				
7	00	SE	83	<u> </u>		LdC	10	)6 3DE)		-		16
6 9	Gd	O7 dE Id	80			Ld HL CALL	LO.	TAPE				17
Ċ	FI	10	80		EXIT	POP	A	F				17
त	87	<del>                                     </del>		<b></b>	-/./	OR	F	7				4
d E	F2	49	81			700	0	ETRAL				
41	B7 F2 AF		<u> </u>	1		XOR	A		,			10
2	19					EXX						4
3	01	04	04			LdBC	0	404				10
3 6 8 9	Ed	79				XOR EXX Ld BC DUT(c) EXX POP	A	-				10 12 4
8	19			<u> </u>		EXX	<u> </u>					4
				<b></b>	RETRN	POP	H	<u> </u>				10
Á	F9		ļ	<u> </u>		Lasp	H	<u>L</u>	<b> </b>			6
B	C9	<b> </b>		ļ	0.55	KET	100	2251	<b> </b>	-14%		10
C	2A	dE 4B	83	<b> </b>	READ	LAHL	18	SPE)	1			10 16 17
F	Cd 18	14B	80	<del>                                     </del>	1	CALL	12	READ XIT	-			12
52 4	00	E8	<b> </b>	<del> </del>	<b>H</b>	JR	F	$\wedge \Box$	<del> </del>			14
T	00	-	<del> </del>	<del> </del>	#	<del> </del>	1		1			T
				1		t			<b>†</b>	<del></del>		
			<u> </u>	<del> </del>			1		1			
				1								
									1			

. , •	. Page 6				-80 VER	SION (	Rev 3	9/22/79	
	IN	STRUCTI	ON			SOURCE COL			
ADDRESS	BYTE 1	BYTE 2	BYTE 3	BYTE 4		OP CODE	OPERAND	COMMENT	\$
·				SER	IAL DA	ATA .	NPUT,	etc.	
8155	78				INPUT	Ld A,	<b>B</b>		4
6	07					RLCA			4
7	d9					EXX			4
8	01		06				0606		10
В		79					, A READR		12.
d	d9					EXX	400-001	0004	
E		F3	83		60.0.T	rg Hr		SCON	16
61	E5	-			SWAIT				
2	IE	FF	<b>}</b>	:	CIONE	Ld E,			7
4		05	<b> </b>		3 LUUN	11/ 2	05 (OS) SERAL		1)
6	<u>d</u> B	05	<b></b>	<del> </del>		RRA ,	US) JENAL		
8	1F	F9	<del>                                     </del>			TO	SLOOK		12/7
9	38		81			CALL	DELAV		77
BE	Cd	75	01			POP	DELAY HL		10
	3E	05	-			Ld A,	05		7
71	-	05				INA.	OS) SERAL		11
3	1F	03	1			RRA	, , , , , , , , , , , , , , , , , , , ,		4
	38	EB		<del>                                     </del>			SWAIT		1217
6	16	FF				Ld d.			7
8		•		1		XUR			4
9	19					EXX			4
À		06	06		:		.0606		10
व	Ed	79					A READR		12
F	19					EXX			4
80	E5				GTBIT	PUSH	HL		11
1	29					ADDITU	HL		11
3	2 B					DEC	1146		6
3	3E	05				Ld A.	05		7
5	9 B	05		·		IN A,	(OF) SERAL		
7	2B 3E dB 5F					LdE	A		4
8	Ca	98	81	1	1	CHLL	DELAY		
B	EI	<u> </u>	ļ	<b></b>	-	POP	HL	<del>                                     </del>	10
C E	38	F2	<u> </u>	<b></b>		JR C	GTBIT	<u> </u>	12/7
E	3E	01	-	<b>}</b>	1	Ld A	0505		4
90		\ <u></u>	1	<u> </u>	1	EXX			10
<u> </u>	01	05	05	-	-}	14 13 L	0505		12
4	Ed	79		<del> </del>		EXX	A SERAL		4
6	49	-	<del> </del>	-{		RET	1	1	10
	69	-	1	<del> </del>	#	KEI			- 70
		-	-	1	1				
	-	-	-	1	#				
	<del>                                     </del>	<b>†</b>	<del>                                     </del>	1			<u> </u>		
	<del>                                     </del>	1	1	1		1			
	1	<del>                                     </del>	-		1				
					1				
				1					
	-			7	i			•	

•	Pa	LGE STRUCTI	7		Z-80	) VE	RSION CO	PE	Rev 3	9/22/79
				,		SOURCE COL	DE			
ADDRESS	BYTE 1	BYTE 2	BYTE 3	BYTE 4		OF CODE	OPERAND	T- 1	COMMENTS	
				54	FRISTL	DELH	y Rou	ואוו		
0105					DCIDY					
8198	7B				DELAY	rg 14,	<u> </u>			4
.9	B0					EXX	13			4
A B	01	05	05				0505			10
E			03			OUT (c)	A SERAL			
	d9	17				EXX	II SERII C			12
1.	28					DEC	141			6
2	7C	-				Ld A,	14			4
ス 3	B5			:		O1S,	L			4
<del>가</del>	20					JR NZ	DELAY	!		12/7
6	7A					Ld A.	d	<u> </u>		4
7	<u>IF</u>					RRA				4
. 8	7B				·	Ld 17,	E			4
9	<u>IF</u>					RRA'				
A	5F					LdE,	1-1			4
B	7A					Ld A.	<u>a</u>			4
2	IF					RRA Ld d,				
d E	57 C9					RET	/ <del>1</del>	<del> </del>		10
E	<u></u>					NE!				
				SER	AL OC	TPOT	ROUTIA	)E		
				JE 1	1110		,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,			
8IAF	87				SROUT	ADD A.	R			4
80	47	,				Ld B,	A			4
J						111 =		200		4
2	3E	OB				Ld A.	11 dec.			<b>ブ</b>
4	4F					rg C;	11 dec.			4
5	17					RLA		ļ		4
6	57					Ld d,	A	ļ		4
<u> </u>	2A	F3	83		OLOOP	rg Hr	A (83F3) SCO (83F3) SCO	<b>Y</b>		16
A	2 A	F3	83			Ld HL	( 83F3) SCO	ν		16
<u> </u>	29				-	Add HI DEC	176			6
=======================================	2B	90	01			COLL	DELAY	ļ		17
4 5 7 7 8 7 8	09 C9	98	81			NEX	DELAY	<u> </u>		4
7.4	20	F2				TR NO	COLOOP	<u> </u>		1217
5	<u>C9</u>					RET				10
<u> </u>	<u>~</u>	<u> </u>	<u> </u>	<b> </b>		1121				
				<b>1</b>						
						ļ		<u> </u>		
			ļ					<u> </u>		
				<b> </b>	<b> </b>	<del> </del>	1	<u> </u>		
		<u></u>	ļ	<b> </b>		1	<b> </b>	<del> </del>		
	<u> </u>	<u> </u>	<b> </b>	<del> </del>	<b> </b>	<b> </b>		<del> </del>		
			<del> </del>	<b></b>		<del>                                     </del>		1	<u> </u>	
<u></u>	<u> </u>	1	<u> </u>	<u>.                                    </u>	11	1	J	<del></del>		

	Pau	1e	8	. 2	-80 VE	ERSION	COPE	Rev 3	9/22/79
ADDRESS	BYTE 1	BYTE 2	BYTE 3	BYTE 4	LABEL	OP CODE		COMMENTS	
		S	ERIL	74	DEVICE	SPEE	D DE	TECTION	
8106	31	EO	83			Ld SP.	83 E O		10
	16	06			DETCT	Lddi	06		7
В	4A					Ld C	4		4
C	OC				TLOOP	Ld C)	<u>C</u>		4
d	21 3E	01	00			Ld HL	0001		10
do	3E	05			GETIN	Ld A	05		
2		05				IN A.C	os) SERAL	1 Timing	here II
4	19			·		EXX	0505	Timing 15 ver	4
5	01	05	05			LE BC,	0505	1 dift ere	10
8	Ed	79				DUT (c),	A SERAL		12 4
A						EXX		<u> </u>	4
В	A9						C		4
Ç	IF					RRA			4 6
d	23					INC	HL		
E	30	FO				JRNL	GETIN		12/7
EO	E5					PUSH	HL		
<u></u>	15					DEC	D TLOOP		4
2	20	E8	ļ			JR N2	TLOOP		12/7
4	EL		<u> </u>			POP	HL		10
5	EL	ļ				POP			10
4 5 6	CI					POP			10
7	09					ADD HL			
8	CI			<u> </u>		POP			10
9 A	09				<b> </b>	ADD HL	BC		
A	CI					POP	BC		0/
B	CI			ļ					10
<u>C</u>	16	03		ļ		Ldd	03		
E	AF				SL001		A		<del></del>
F	7C			<u> </u>		Ld A	11		<del></del>
Fo	IF	ļ		<u> </u>	<b> </b>	RRA			
	67	<b></b>	ļ		<u> </u>	Ld 14,	I P		4
2 3	7d	<u> </u>	<u> </u>			Ld It,	<u> </u>		
	IF		ļ	<b>_</b>		RRA			7
4	6F		ļ	<del> </del>	1	Ld L	119	<u> </u>	
5	15	<del> </del>		<b>_</b>	<del>  </del>	DEC			4
6		F6	ļ	1	-	JK NZ	SLOOP		147
8	2B	<del> </del>		<u> </u>	-	DEC	116		16
9	22	F3	83	<del>- </del>	H	Trg (831	HL B),HL Scon STOP	4	
C	18	FE		<b>_</b>	STOP	1111	SIOP		12
E	00	00	-		#		+		
	ļ	<b>-</b>	- <del> </del>						
		<del> </del>	1		#	1 .			
	<del> </del>	<del> </del>		-		BERT F. HASSAR ice Creek Drive		·	
	ļ	<del> </del>	-	1		r Creek, Calif. 9			
	-	-		<del> </del>	#				
	<del> </del>	+	+	1	#		<del> </del>		
	<del> </del>	-	+		-				
L					<del>-!l</del>			i	